

## Conservative Entailment

Gert Smolka, May 5, 2005

## Solutions

$\sigma$  is solution of  $E$  in  $\mathcal{D}$  if

$$\forall s=t \in E. \mathcal{D}\sigma = \mathcal{D}\sigma t$$

$$\text{Sol}_{\mathcal{D}} E \stackrel{\text{def}}{=} \{ \sigma \mid \sigma \text{ solution of } E \text{ in } \mathcal{D} \}$$

$$\text{Sol}_{\mathcal{D}}(E_1 \cup E_2) = \text{Sol}_{\mathcal{D}} E_1 \cap \text{Sol}_{\mathcal{D}} E_2$$

## Solving Equation Systems

Solving  $E$  means to derive  $E'$  p.t.

1)  $\text{Sol}_{\mathcal{D}} E = \text{Sol}_{\mathcal{D}} E'$

2) Solutions of  $E'$  are obvious

$E'$  is called solved form of  $E$

## Example

$$\begin{cases} x+2y=1 \\ x-y=1 \end{cases} \xrightarrow{\mathcal{Z}} \begin{cases} x+2y=1 \\ x=y+1 \end{cases}$$

$$\xrightarrow{\mathcal{Z}} \begin{cases} (y+1)+2y=1 \\ x=y+1 \end{cases} \xrightarrow{\mathcal{Z}} \begin{cases} 3y=0 \\ x=y+1 \end{cases}$$

$$\xrightarrow{\mathcal{Z}} \begin{cases} x=y+1 \\ y=0 \end{cases} \xrightarrow{\mathcal{Z}} \begin{cases} x=0+1 \\ y=0 \end{cases} \xrightarrow{\mathcal{Z}} \begin{cases} x=1 \\ y=0 \end{cases}$$

## Conservative Entailment

$$E \stackrel{\mathcal{D}}{\vdash}_0 e \stackrel{\text{def}}{\iff} \text{Sol}_{\mathcal{D}} E \subseteq \text{Sol}_{\mathcal{D}} \{e\}$$

$$E \stackrel{\mathcal{D}}{\vdash}_0 e \implies \exists \theta E \stackrel{\mathcal{D}}{\vdash} \theta e$$

stable under substitution

## Conservative Deduction

$$E \vdash_0 e \stackrel{\text{def}}{\iff} E \vdash e \text{ without Subst and capture}$$

$$E \vdash^A_0 e \stackrel{\text{def}}{\iff} E \cup \text{DCA} \vdash_0 e$$

$$\mathcal{D} \models A \wedge E \vdash^A_0 e \implies E \stackrel{\mathcal{D}}{\vdash}_0 e$$

So on these

$$E \vdash^A_0 e \implies \exists \theta E \vdash^A_0 \theta e$$

Stability under substitution

## Conservative Rewriting

$$\Delta \stackrel{E}{\rightrightarrows}_0 t \stackrel{\text{def}}{\iff} \Delta \stackrel{E}{\rightrightarrows} t \text{ without Subst and without capturing}$$

$$E \vdash^A_0 s = t \iff s = t \vee s (\stackrel{E}{\rightrightarrows}_0 \cup \stackrel{A}{\rightrightarrows})^+ t$$

## Capture-free Replacement

$$E_1 \vdash^E_0 E_2 \iff E_1 \stackrel{E}{\vdash}_0 E_2 \wedge E_2 \stackrel{E}{\vdash}_0 E_1$$

$$E \cup \{x=y\} \vdash_0 [E[x:=y]] \cup \{x=y\}$$