



Assignment 10

Introduction to Computational Logic, SS 2011

Prof. Dr. Gert Smolka, Dr. Chad Brown

www.ps.uni-saarland.de/courses/cl-ss11/

Read in the lecture notes: Chapters 5-6

Note: It is very important to do all the examples in the lecture notes and the exercises below in the system Coq.

Exercise 10.1 Prove the following lemmas.

Lemma ex5811 $n : \text{even } n \rightarrow \text{even } (4+n)$.

Lemma ex5812 $n : \text{even } (S (S n)) \rightarrow \text{even } n$.

Exercise 10.2 Prove the following lemmas.

Lemma ex58221 : $\sim \text{even } 3$.

Lemma ex5822 $n : \text{even } (4+n) \rightarrow \text{even } n$.

Exercise 10.3 Prove the following lemmas. Do not use induction on *nat*; only use induction on proofs of propositions of the form *evens*.

Lemma even_sum $m n : \text{even } m \rightarrow \text{even } n \rightarrow \text{even } (m+n)$.

Lemma even_sum' $m n : \text{even } (m+n) \rightarrow \text{even } m \rightarrow \text{even } n$.

Exercise 10.4 Prove that the inductive and the boolean definitions of evenness are equivalent.

Lemma evenib $n : \text{even } n \leftrightarrow \text{evenb } n$.

Exercise 10.5 Here is an impredicative definition of evenness.

Definition evenp $(n : \text{nat}) : \text{Prop} :=$

forall $p : \text{nat} \rightarrow \text{Prop}$,

$p 0 \rightarrow (\text{forall } n, p n \rightarrow p (S (S n))) \rightarrow p n$.

Prove that the inductive and the impredicative definitions of evenness are equivalent.

Lemma evenip $n : \text{even } n \leftrightarrow \text{evenp } n$.

Exercise 10.6 Consider the following inductive definition of equality where the type constructor $eq2$ takes two proper arguments.

Inductive $eq2$ ($X : Type$) : $X \rightarrow X \rightarrow Prop$:=
 | $eq2_I$: $\text{forall } x : X, eq2\ X\ x\ x$.

- Give the typing rule for matches at $eq2$.
- Prove the following property of $eq2$.

Lemma $eq2_E$ ($X : Type$) ($x\ y : X$) :
 $eq2\ X\ x\ y \rightarrow \text{forall } p : X \rightarrow Prop, p\ y \rightarrow p\ x$.

Exercise 10.7 Consider the following inductive definition of an order predicate for the natural numbers.

Inductive lei : $nat \rightarrow nat \rightarrow Prop$:=
 | $leiO$: $\text{forall } x : nat, lei\ 0\ x$
 | $leiS$: $\text{forall } x\ y, lei\ x\ y \rightarrow lei\ (S\ x)\ (S\ y)$.

- Given the typing rule for matches at lei .
- Prove the following lemmas.

Lemma lei_refl $x : lei\ x\ x$.
Lemma lei_trans $x\ y\ z : lei\ x\ y \rightarrow lei\ y\ z \rightarrow lei\ x\ z$.
Lemma $leib$ $x\ y : leb\ x\ y \leftrightarrow lei\ x\ y$.

Exercise 10.8 Use proof scripts to give inhabitants of the following two types.

- $\forall p : Prop. p \vee p \rightarrow p + p$
- $\forall p : Prop. p \vee False \rightarrow p + False$

Exercise 10.9 Complete the following definition and prove the lemma.

Definition $\text{forget } \{X\ Y\} \{p : X \rightarrow Prop\} \{q : Y \rightarrow Prop\} : sig\ p + sig\ q \rightarrow X + Y$.
Lemma forget_div2c ($n : nat$) : $\text{forget } (div2c\ n) = divmod2\ n$.

Exercise 10.10 Recall the definition of the type $Search$ from the lecture.

Definition $Search : Type$:=
 $\text{forall } (p : nat \rightarrow bool) (n : nat),$
 $\{x \mid x \leq n \wedge p\ x\} + (\text{forall } x, x \leq n \rightarrow \sim p\ x)$.

Define a similar type $SearchMax$ such that any function of type $SearchMax$ will return the maximum $x \leq n$ such that $p\ x$ if such an x exists, or a proof that no such x exists. Construct a certifying function of this type.