



## Semantics of Programming Languages Assignment 12

Andreas Podelski, Patrick Maier, Jan Schwinghammer

<http://www.ps.uni-sb.de/courses/sem-ws01/>



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Consider the following program  $P$ ; the integer variables  $x$  and  $y$  both are initialized to 0.

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11: while (*) {
12:   x++;
13:   y++;
14: }
15: while (*) {
16:   x--;
17:   y--;
18: }
19: if (x == y) {
20:   assert(0);
21: }
```

**Exercise 12.1 (10)** Let the set of predicates be  $\{p_1\}$  where  $p_1 \equiv x = y$ .

1. For each of the assignments in  $P$ , compute the corresponding abstract statement.
2. Construct the abstract program  $P^\#$  such that  $post$  of  $P^\#$  equals  $post_{cart}^\#$  of  $P$ .
3. Compute the collecting semantics of  $P^\#$ . Is the label  $l8$  reachable in  $P^\#$ ?

**Exercise 12.2 (5)** Refine the abstraction (i. e., specify additional predicates) from exercise 12.1 such that the label  $l8$  is not reachable in the abstract program. Give the new abstract program (without proof).

**Exercise 12.3 (5)** Formalize constant propagation (see exercise 9.3) as predicate abstraction, i. e., give the predicates and the boolean abstraction.

Why is this analysis not implementable? Suggest a way to make it implementable, at least for some programs.

Discuss the effect of a cartesian abstraction (see exercises 8.2 and 8.3) on top of the boolean abstraction.

**Exercise 12.4 (5)** Let  $\mathbf{D} = \langle D, \sqsubseteq \rangle$  and  $\mathbf{D}^\# = \langle D^\#, \sqsubseteq^\# \rangle$  be posets. A function  $f : D \rightarrow D^\#$  is called *additive* if for all  $a, b \in D$ , the existence of  $a \sqcup b$  implies that  $f(a) \sqcup^\# f(b)$  exists and  $f(a \sqcup b) = f(a) \sqcup^\# f(b)$ .

Let  $\mathbf{D} \xrightarrow[\alpha]{\gamma} \mathbf{D}^\#$  be a Galois connection. Prove that  $\alpha$  is additive. Give a counterexample showing that  $\gamma$  need not be additive even if  $\gamma$  is the identity, i. e.,  $\gamma(y) = y$  for all  $y \in D^\#$ .

Generalize the concept of additive functions to *completely additive* functions and prove that  $\alpha$  is completely additive. What related property holds for  $\gamma$ ?

Just for fun: Considering the posets  $\mathbf{D}$  and  $\mathbf{D}^\#$  as CPOs, what are the relations between additive, completely additive and continuous functions?