



## Semantics, WS 2003 – Assignment 3

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<http://www.ps.uni-sb.de/courses/sem-ws03/>

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Recommended reading: Types and programming languages, chapters 8,9

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**Exercise 3.1: Reference sheet I** Write a reference sheet (one page only) for the simply typed  $\lambda$ -calculus with Bool and Nat containing the definitions of

- (a) Syntax:  $T \in \text{Typ}, x \in \text{Var}, t \in \text{Ter}, v \in \text{Val}, nv \in \text{NVal}$
- (b) Reduction:  $t \rightarrow t'$
- (c) Typing:  $\Gamma \in \text{TE}, \Gamma \vdash t : T$

Bring this reference sheet to the lecture and use it for all proofs.

**Exercise 3.2: Reference sheet II** Write a reference sheet (one page only) for the simply typed  $\lambda$ -calculus with Bool and Nat containing the formulation of the following properties:

- (a) Uniqueness
- (b) Progress
- (c) Substitution
- (d) Preservation
- (e) Normalization

Remark for each of the properties by which induction they can be proved.

**Exercise 3.3: Big step semantics** The so-called *big step semantics* is defined as follows:

$$t \Downarrow v :\Leftrightarrow t \rightarrow^* v$$

Define the big step semantics of the  $\lambda$ -calculus with Bool and Nat independently of  $\rightarrow$  by inference rules. Exercise 3.5.17 in the book can give you some hints.

**Exercise 3.4: Interpreter for the simply typed  $\lambda$ -calculus** We extend the syntax of terms by boolean constants, an **IF** expression and types in the following way:

```
datatype typ = ARROW of typ * typ
             | BOOL
             | INT

type var = int

datatype term = V of var
              | L of typ * term
              | A of term * term
              | FALSE
              | TRUE
              | IF of term * term * term
              | ZERO
              | SUCC
              | PRED
              | ISZERO
```

Type environments are implemented as a function  $\gamma : \text{var} \rightarrow \text{typ}$ .

- Write a procedure  $\text{typeof} : \text{term} \rightarrow \text{typ}$  that returns the type of a term or raises an exception if the term has no type.
- Adjust the procedures  $\text{shift}$  and  $\text{subst}$  from the previous assignment sheet to the extended term datatype.
- Write a procedure  $\text{reduce} : \text{term} \rightarrow \text{term}$  that yields the normal form of a closed term.

**Exercise 3.5: Induction Theorem** Write the precise formulation of the Induction Theorem. Try to prove it.

**Exercise 3.6: Rule induction** Write the precise definitions of  $R$ ,  $\hat{R}$ , and  $I_R$  for ground rule systems. Give the precise formulation of the Rule Induction Theorem. Try to prove it by induction on derivations.

The set  $I_R$  can be characterized as the least subset of  $X$  satisfying a certain property. State this property in full detail.

**Exercise 3.7: Substitution Lemma** Here is a statement of the Substitution Lemma:

$$\Gamma[x : S] \vdash t : T \wedge \Gamma \vdash s : S \implies \Gamma \vdash t[x := s] : T$$

The variables  $\Gamma, x, s, t, T, S$  are all universally quantified over their canonical domains.

The lemma can be proved by induction over the rules defining the typing relation (see the book, page 106). State the set  $P$  that the proof is using (with respect to the Rule Induction Proposition). Make sure that you describe  $P$  in full detail but as concisely as possible.

**Exercise 3.8: Reversed Type Preservation** Prove the following statement:

$$\neg(t \rightarrow t' \wedge \Gamma \vdash t' : T) \implies \Gamma \vdash t : T$$

Intuitively, the statement says that type preservation doesn't hold from right to left.