



## Assignment 1 Semantics, WS 2007/08

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Standard ML: recursion, cascaded procedures

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**Exercise 1.1 (Fib from iter)** Write a procedure  $fib : int \rightarrow int$  such that

$$fib\ n = \text{if } n < 2 \text{ then } n \text{ else } fib(n - 1) + fib(n - 2)$$

Use the procedure

```
fun iter n x f = if n < 1 then x else iter (n-1) (f x) f
```

to obtain the necessary recursion.

**Exercise 1.2 (Gcd from loop)** Consider the procedure

```
fun loop x p f = if p x then loop (f x) p f else x
```

Write a procedure  $gcd : int \rightarrow int \rightarrow int$  that yields the greatest common divisor for positive arguments. Use the equation

$$gcd\ x\ y = \text{if } x = 0 \text{ then } y \text{ else } gcd(y \bmod x)\ x$$

and use *loop* to realize the necessary recursion.

**Exercise 1.3 (Mapr from foldl)** Write a procedure  $mapr : (\alpha \rightarrow \beta) \rightarrow \alpha\ list \rightarrow \beta\ list$  such that

$$mapr\ f\ [x_1, \dots, x_n] = [f\ x_n, \dots, f\ x_1]$$

To obtain the necessary recursion, use the following procedure:

```
fun foldl f s nil = s
  | foldl f s (x::xr) = foldl f (f(x,s)) xr
```

**Exercise 1.4 (List from iter)** Write a procedure  $list : int \rightarrow (int \rightarrow \alpha) \rightarrow \alpha\ list$  such that

$$list\ n\ f = [f\ 1, \dots, f\ n]$$

using *iter* instead of a recursive definition.

**Exercise 1.5 (Power from fix)** Write a procedure  $power : int \rightarrow int \rightarrow int$  such that  $power\ x\ n = x^n$ . Use *fix* to realize the necessary recursion, where

```
fun fix f x = f (fix f) x
```

**Exercise 1.6 (Cas and car)** Write procedures

$$cas : (\alpha * \beta \rightarrow \gamma) \rightarrow \alpha \rightarrow \beta \rightarrow \gamma$$
$$car : (\alpha \rightarrow \beta \rightarrow \gamma) \rightarrow \alpha * \beta \rightarrow \gamma$$

such that  $car(cas\ f) = f$  and  $cas(car\ g) = g$ .

**Exercise 1.7 (Cond)** Write a procedure  $cond : bool \rightarrow (unit \rightarrow \alpha) \rightarrow (unit \rightarrow \alpha) \rightarrow \alpha$  such that

$$cond\ t_0\ (fn\ () \Rightarrow t_1)\ (fn\ () \Rightarrow t_2)$$

is semantically equivalent to the expression *if*  $t_0$  *then*  $t_1$  *else*  $t_2$ .