



## Assignment 5 Semantics, WS 2007/08

Prof. Dr. Gert Smolka, Dr. Jan Schwinghammer  
[www.ps.uni-sb.de/courses/sem-ws07/](http://www.ps.uni-sb.de/courses/sem-ws07/)

---

Read Chapter 3 of the lecture notes

---

**Exercise 5.1 (Substitution with  $\rho_0$ )** Determine  $S_0 \in t$  for the following terms  $t$ . Assume  $x = 5, y = 0$  and  $z = 3$ .

- a)  $\lambda x.x$ ,
- b)  $\lambda xy.x(x y)$ ,
- c)  $\lambda xyz.y(x z)$ ,
- d)  $\lambda yzx.z(y x)$ ,
- e)  $(\lambda x.y) y$ .

**Exercise 5.2 ( $\alpha$ -renaming)** Find counterexamples that falsify the following statements:

- a)  $\lambda x.s \sim_\alpha \lambda y.t \iff \exists z : [x := z]s \sim_\alpha [y := z]t$
- b)  $\lambda x.s \sim_\alpha \lambda y.t \iff t \sim_\alpha [x := y]s$

**Exercise 5.3 (Big step semantics)** Define  $\Downarrow \subseteq Ter \times Val$  with inference rules for the untyped call-by-value  $\lambda$  calculus. Give two versions, one for  $Val = Val_1$  and one for  $Val = Val_2$ . Make sure that for all  $t$  and  $v, t \Downarrow v \iff t \rightarrow \dots \rightarrow v$  holds for both versions.

**Exercise 5.4 (Higher-order abstract syntax)** In this exercise you will implement the untyped call-by-value  $\lambda$  calculus with  $Val_2$  in Standard ML using higher-order abstract syntax (HOAS). Use the following declarations:

```
structure Var :>
sig
  eqtype var
  val var : unit -> var
end =
struct
  type var = int
  val r = ref ~1
  fun var () = (r := !r + 1; !r)
end

open Var
datatype ter = V of var | A of ter * ter | L of ter -> ter
```

- a) Write a procedure  $eval : ter \rightarrow ter$  that evaluates terms.
- b) Write procedures  $church : int \rightarrow ter$  and  $dechurch : ter \rightarrow int$  such that  $church\ n$  yields a value equivalent to  $c_n$ , and  $dechurch\ t$  yields  $n$  if  $t \sim c_n$ .
- c) Give a term  $power : ter$  such that  $power\ c_m\ c_n$  yields  $c_{m^n}$ .
- d) Write a procedure  $size : ter \rightarrow int$  that yields the size of a term.
- e) Write a procedure  $alpha : ter \rightarrow ter \rightarrow bool$  that tests whether two terms are  $\alpha$ -equivalent.
- f) Write a procedure  $free : var \rightarrow ter \rightarrow bool$  that tests if a variable occurs free in a term.
- g) Write a procedure  $subst : (var \rightarrow ter) \rightarrow ter \rightarrow ter$  that applies a substitution to a term.