



## Assignment 6 Semantics, WS 2007/08

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Read Chapter 3 of the lecture notes and Chapters 6, 8 and 9 of Pierce's TAPL

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### Exercise 6.1 (Contextual equivalence)

- Show that  $s \sim t \implies \theta s \sim \theta t$  doesn't hold in the pure calculus.
- Show that  $x \sim (\lambda y.x)y$  doesn't hold if contextual equivalence is defined without the closedness condition.
- Show that  $x \approx \lambda y.x y$  doesn't hold in the extended calculus.

**Exercise 6.2 (De Bruijn terms)** We implement ordinary and de Bruijn terms as follows:

```
type var = int
datatype ter = V of var | A of ter * ter | L of var * ter
datatype dbt = DV of var | DA of dbt * dbt | DL of dbt
```

- Write a procedure  $free : var \rightarrow dbt \rightarrow bool$  that tests whether a variable occurs free in a de Bruijn term.
- Write a procedure  $subst : (var \rightarrow dbt) \rightarrow dbt \rightarrow dbt$  that applies a de Bruijn substitution to a de Bruijn term.
- Write a procedure  $db : ter \rightarrow dbt$  that translates a term into a de Bruijn term.

### Exercise 6.3 (Simple types)

- Show that for any  $T$ , the  $n$ -th Church numeral  $c_n := \lambda f:T \rightarrow T. \lambda x:T. f^n x$  has type  $(T \rightarrow T) \rightarrow T \rightarrow T$ .
- Give a type-annotated version of  $succ$  (see Sect. 3.4 of the notes), and verify that it has type  $((T \rightarrow T) \rightarrow T \rightarrow T) \rightarrow (T \rightarrow T) \rightarrow T \rightarrow T$ . Can you do the same for  $power = \lambda mn.n m$ , taking 2 arguments of type  $(T \rightarrow T) \rightarrow T \rightarrow T$ ?
- Show that  $\omega = \lambda x:T.x x$  is not well-typed, for any choice of  $T$ . (By Exercise 6.6 below, you may therefore conclude that also  $\Omega = \omega \omega$  is not typable.) **Hint:** use the Inversion Lemma.

**Exercise 6.4 (Expansion)** Show that *expansion* does not hold: find  $\Gamma, t, t'$  and  $T$  such that  $\Gamma \vdash t' : T$  and  $t \rightarrow t'$ , but  $\Gamma \vdash t : T$  does not hold.

**Exercise 6.5 (Relevance)** Prove that  $\Gamma \vdash t : T$  and  $x \in FV t$  implies that  $x \in Dom \Gamma$ .