



## Assignment 7 Semantics, WS 2011-2012

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[www.ps.uni-saarland.de/courses/c1-ss11/](http://www.ps.uni-saarland.de/courses/c1-ss11/)

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Read in the lecture notes:

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Read the new material in Chapter 4 of the lecture notes.

**Exercise 7.1** Give the induction principles for the following inductive predicates defined in this chapter and check your results with Coq.

- a) *And*
- b) *Eq*
- c) *EQ*

**Exercise 7.2** Give the induction principle for the following inductive predicate.

**Inductive** `le2` : `nat`  $\rightarrow$  `nat`  $\rightarrow$  `Prop` :=  
| `le2x` : `forall` `x`, `le2` `x` `x`  
| `le2S` : `forall` `x` `y`, `le2` `x` `y`  $\rightarrow$  `le2` `x` (`S` `y`).

**Exercise 7.3** Prove the following goal. Do not use a lemma.

**Lemma** `starTR` : `forall` `x` `y` `z`, `star` `x` `y`  $\rightarrow$  `R` `y` `z`  $\rightarrow$  `star` `x` `z`.

**Exercise 7.4** Give the induction principle for the inductive predicate *star*.

**Exercise 7.5** We can define a reflexive transitive closure predicate *star1* with a single proper argument.

**Inductive** `star1` (`x` : `X`) : `X`  $\rightarrow$  `Prop` :=  
| `star1R` : `star1` `x` `x`  
| `star1T` : `forall` `y` `z`, `star1` `x` `y`  $\rightarrow$  `R` `y` `z`  $\rightarrow$  `star1` `x` `z`.

- a) Give the induction principle for *star1*.
- b) Prove that *star1* is reflexive and transitive.
- c) Prove  $\forall xy. \text{star } xy \leftrightarrow \text{star1 } xy$

**Exercise 7.6** Prove that taking the reflexive transitive closure preserves invariants.

**Definition** `invariant` {`X` : `Type`} (`p` : `X`  $\rightarrow$  `Prop`) (`R` : `X`  $\rightarrow$  `X`  $\rightarrow$  `Prop`) : `Prop` :=  
`forall` `x` `y`, `R` `x` `y`  $\rightarrow$  `p` `x`  $\rightarrow$  `p` `y`.

Goal forall (X : Type) (R : X -> X -> Prop) (p : X -> Prop),  
invariant p R -> invariant p (star R).

**Exercise 7.7** You may have seen  $R^* := \bigcup_{n \in \mathbb{N}} R^n$  as a definition of the reflexive transitive closure. Using the function *iter*, we can express this definition in Coq.

**Definition** comp {X : Type} (R S : X -> X -> Prop) (x z : X) : Prop :=  
exists y, R x y /\ S y z.

**Definition** stari {X : Type} (R : X -> X -> Prop) (x y : X) :=  
exists n, iter n (comp R) (fun x y => x=y) x y.

Prove the equivalence of the inductive and the iterative definition.

Goal forall (X : Type) (R : X -> X -> Prop) (x y : X),  
star R x y <-> stari R x y.

**Exercise 7.8** Give three proofs for the following goal.

Goal forall r r' : rel, rap r r' -> functional r' -> functional r.

- a) Use *firstorder*.
- b) Use *eauto*.
- c) Use *eapply* and *eassumption*.

**Exercise 7.9** Download the Coq code from the lecture of November 30. Consider the correctness theorem for the compiler from commands to regular expressions.

**Theorem** correctness c :  
req (sem c) (den (compile c)).

The proof has two directions. The proof of the first direction is given.

- a) Rewrite the proof of the first direction so that it does not use *firstorder* or *eauto*.
- b) Write the proof of the second direction.

**Exercise 7.10** Prove the following goal by applying the induction principle *even\_ind*. Do not use the induction tactic.

Goal forall n, even n -> even (S n) -> False.