



Semantics, WS 2011-2012: Solution for Assignment 7

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Read the new material in Chapter 4 of the lecture notes.

Exercise 7.1 Give the induction principles for the following inductive predicates defined in this chapter and check your results with Coq.

- a) *And*
- b) *Eq*
- c) *EQ*

Solution to Exercise 7.1

Inductive *And* (X Y : Prop) : Prop :=
| *AndI* : X → Y → And X Y.

And_ind

: forall X Y P : Prop, (X → Y → P) → And X Y → P

Inductive *Eq* (X : Type) : X → X → Prop :=
| *EqI* : forall x, Eq X x x.

Eq_ind

: forall (X : Type) (P : X → X → Prop),
(forall x : X, P x x) → forall y z : X, Eq X y z → P y z

Inductive *EQ* (X : Type) (x : X) : X → Prop :=
| *EQI* : EQ X x x.

EQ_ind

: forall (X : Type) (x : X) (P : X → Prop),
P x → forall y : X, EQ X x y → P y

Exercise 7.2 Give the induction principle for the following inductive predicate.

Inductive *le2* : nat → nat → Prop :=
| *le2x* : forall x, le2 x x
| *le2S* : forall x y, le2 x y → le2 x (S y).

Solution to Exercise 7.2

Inductive $\text{le2} : \text{nat} \rightarrow \text{nat} \rightarrow \text{Prop} :=$
| $\text{le2x} : \text{forall } x, \text{le2 } x \ x$
| $\text{le2S} : \text{forall } x \ y, \text{le2 } x \ y \rightarrow \text{le2 } x \ (S \ y).$

le2_ind
: $\text{forall } P : \text{nat} \rightarrow \text{nat} \rightarrow \text{Prop},$
 $(\text{forall } x : \text{nat}, P \ x \ x) \rightarrow$
 $(\text{forall } x \ y : \text{nat}, \text{le2 } x \ y \rightarrow P \ x \ y \rightarrow P \ x \ (S \ y)) \rightarrow$
 $\text{forall } x \ y : \text{nat}, \text{le2 } x \ y \rightarrow P \ x \ y$

Exercise 7.3 Prove the following goal. Do not use a lemma.

Lemma $\text{starTR} : \text{forall } x \ y \ z, \text{star } x \ y \rightarrow R \ y \ z \rightarrow \text{star } x \ z.$

Solution to Exercise 7.3

Lemma $\text{starTR} : \text{forall } x \ y \ z, \text{star } x \ y \rightarrow R \ y \ z \rightarrow \text{star } x \ z.$

Proof. $\text{intros } x \ y \ z \ A \ B.$ induction A ; eauto using star. **Qed.**

Exercise 7.4 Give the induction principle for the inductive predicate *star*.

Solution to Exercise 7.4

star_ind
: $\text{forall } P : X \rightarrow X \rightarrow \text{Prop},$
 $(\text{forall } x : X, P \ x \ x) \rightarrow$
 $(\text{forall } x \ y \ z : X, R \ x \ y \rightarrow \text{star } y \ z \rightarrow P \ y \ z \rightarrow P \ x \ z) \rightarrow$
 $\text{forall } x \ y : X, \text{star } x \ y \rightarrow P \ x \ y$

Exercise 7.5 We can define a reflexive transitive closure predicate *star1* with a single proper argument.

Inductive $\text{star1} (x : X) : X \rightarrow \text{Prop} :=$
| $\text{star1R} : \text{star1 } x \ x$
| $\text{star1T} : \text{forall } y \ z, \text{star1 } x \ y \rightarrow R \ y \ z \rightarrow \text{star1 } x \ z.$

- Give the induction principle for *star1*.
- Prove that *star1* is reflexive and transitive.
- Prove $\forall xy. \text{star } xy \leftrightarrow \text{star1 } xy$

Solution to Exercise 7.5

```
star1_ind
: forall (x : X) (P : X -> Prop),
  P x ->
  (forall y z : X, star1 x y -> P y -> R y z -> P z) ->
  forall y : X, star1 x y -> P y
```

Goal forall x, star1 x x.

Proof. eauto using star1. **Qed.**

Goal forall x y z, star1 x y -> star1 y z -> star1 x z.

Proof. intros x y z A B. induction B ; eauto using star1. **Qed.**

Lemma star1TL : forall x y z, R x y -> star1 y z -> star1 x z.

Proof. intros x y z A B. induction B ; eauto using star1. **Qed.**

Goal forall x y, star x y <-> star1 x y.

Proof. split.

intros A. induction A ; eauto using star1, star1TL.

intros A. induction A ; eauto using star, starTR.

Qed.

Exercise 7.6 Prove that taking the reflexive transitive closure preserves invariants.

Definition invariant {X : Type} (p : X -> Prop) (R : X -> X -> Prop) : Prop :=
forall x y, R x y -> p x -> p y.

Goal forall (X : Type) (R : X -> X -> Prop) (p : X -> Prop),
invariant p R -> invariant p (star R).

Solution to Exercise 7.6

Goal forall (X : Type) (R : X -> X -> Prop) (p : X -> Prop),
invariant p R -> invariant p (star R).

Proof. intros X r p A x y C D. induction C ; firstorder. **Qed.**

Exercise 7.7 You may have seen $R^* := \bigcup_{n \in \mathbb{N}} R^n$ as a definition of the reflexive transitive closure. Using the function *iter*, we can express this definition in Coq.

Definition `comp` $\{X : \text{Type}\} (R S : X \rightarrow X \rightarrow \text{Prop}) (x z : X) : \text{Prop} :=$
`exists y, R x y /\ S y z.`

Definition `stari` $\{X : \text{Type}\} (R : X \rightarrow X \rightarrow \text{Prop}) (x y : X) :=$
`exists n, iter n (comp R) (fun x y => x=y) x y.`

Prove the equivalence of the inductive and the iterative definition.

Goal `forall (X : Type) (R : X -> X-> Prop) (x y : X),`
`star R x y <-> stari R x y.`

Solution to Exercise 7.7

Goal `forall (X : Type) (R : X -> X-> Prop) (x y : X),`
`star R x y <-> stari R x y.`

Proof. `split.`

`intros A. induction A. exists 0. reflexivity.`

`destruct IHA as [n B]. exists (S n). simpl. exists y. tauto.`

`intros [n A]. revert x y A. induction n ;simpl.`

`intros x y A. rewrite A. now constructor.`

`intros x y [x' [A B]]. apply (starT X R x x') ; now auto. Qed.`

Exercise 7.8 Give three proofs for the following goal.

Goal `forall r r' : rel, rap r r' -> functional r' -> functional r.`

- a) Use *firstorder*.
- b) Use *eauto*.
- c) Use *eapply* and *eassumption*.

Solution to Exercise 7.8

Goal `forall r r' : rel, rap r r' -> functional r' -> functional r.`

Proof. `firstorder. Qed.`

Goal `forall r r' : rel, rap r r' -> functional r' -> functional r.`

Proof. `intros r r' A B s s' s'' C D. eauto. Qed.`

Goal `forall r r' : rel, rap r r' -> functional r' -> functional r.`

Proof. `intros r r' A B s s' s'' C D. eapply B; eapply A; eassumption. Qed.`

Exercise 7.9 Download the Coq code from the lecture of November 30. Consider the correctness theorem for the compiler from commands to regular expressions.

Theorem correctness c :
req (sem c) (den (compile c)).

The proof has two directions. The proof of the first direction is given.

- a) Rewrite the proof of the first direction so that it does not use *firstorder* or *eauto*.
- b) Write the proof of the second direction.

Solution to Exercise 7.9

Theorem correctness c :
req (sem c) (den (compile c)).

Proof. split.

```
(* -> *)
intros s s' A. induction A.
(* Act *) simpl. reflexivity.
(* Seq *) simpl. exists s'. split; assumption.
(* IfTrue *) simpl. left. exists s. split. split.
      assumption. reflexivity. assumption.
(* IfFalse *) simpl. right. exists s. split. split.
      apply negteq. assumption. reflexivity. assumption.
(* WhileFalse *) simpl. exists s. split. now apply starR.
      split. apply negteq. assumption. reflexivity.
(* WhileTrue *) simpl. exists s''.
      destruct IHA2 as [s3 [H3 [H4 H5]]]. subst.
      split.
      apply starS with (s' := s').
      exists s. split. split. assumption. reflexivity. assumption.
      assumption.
      split. assumption. reflexivity.
(* <- *)
induction c; simpl; intros s s' A.
(* Act *) rewrite <- A. now constructor.
(* Seq *) destruct A as [s2 [H1 H2]]. econstructor. eapply IHc1. eassumption.
      apply IHc2. now apply H2.
(* If *) destruct A as [[s2 [[H1 H2] H3]][s2 [[H1 H2] H3]]].
      subst. econstructor. assumption. apply IHc1. assumption.
      subst. apply semIfFalse. apply negteq. assumption. apply IHc2. assumption.
```

```
(* While *) destruct A as [s2 [H [H1 H2]]]. subst.  
  induction H.  
  apply semWhileFalse. apply negteq. assumption.  
  destruct H as [s3 [[H2 H3] H4]]. subst.  
  apply semWhileTrue with (s' := s'). assumption. apply IHc. assumption.  
  apply IHstar. assumption.
```

Qed.

Exercise 7.10 Prove the following goal by applying the induction principle *even_ind*. Do not use the induction tactic.

Goal forall n, even n -> even (S n) -> False.

Solution to Exercise 7.10

```
Goal forall n, even n -> even (S n) -> False.  
apply (even_ind (fun n => even (S n) -> False)).  
intros A. now inversion A.  
intros n A IHA B. inversion B. auto. Qed.
```