



Assignment 3

Introduction to Computational Logic, SS 2008

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Read in the lecture notes: Chapters 3 and 4

Exercise 3.1 Give finite relations \rightarrow such that:

- a) \rightarrow is confluent but not terminating.
- b) \rightarrow is terminating but not confluent.
- c) \rightarrow is not confluent and not terminating.
- d) \rightarrow is confluent, does not terminate on x , and y is a \rightarrow -normal form of x .

Exercise 3.2 Consider the relation $\rightarrow := \{ (x, y) \in \mathbb{N}^2 \mid 2 \leq 2y \leq x \}$.

- a) Is \rightarrow terminating?
- b) Is \rightarrow confluent?
- c) Give a \rightarrow -normal form of 7.
- d) Give all \rightarrow -normal $n \in \mathbb{N}$.

Exercise 3.3 A relation \rightarrow is *locally confluent* if for all x, y, z : $x \rightarrow y \wedge x \rightarrow z \implies y \downarrow z$. Find a finite relation that is locally confluent but not confluent.

Exercise 3.4 Find terms s, t such that $s \rightarrow_\beta t$, s contains no η -redex, and t contains an η -redex.

Exercise 3.5 Which condition in the definition of interpretations ensures that sorts are interpreted as non-empty sets?

Exercise 3.6 (Multiplication) Extend the specification of the natural number with a formula that specifies the name $\cdot : NNN$ as multiplication.

Exercise 3.7 (Pairs) Let the names $\text{pair} : \sigma\tau P$, $\text{fst} : P\sigma$, and $\text{snd} : P\tau$ be given. Find a formula that is satisfied by a logical interpretation \mathcal{I} if and only if $\mathcal{I}P \cong \mathcal{I}\sigma \times \mathcal{I}\tau$ and pair , fst , and snd are interpreted as the pairing and projection functions.

Exercise 3.8 (Termination) Let $r : \alpha\alpha B$ be a name. Find a formula that is satisfied by a logical interpretation \mathcal{I} if and only if $\mathcal{I}r$ is the functional coding of a terminating relation.

Exercise 3.9 (Finiteness) Let $f : \sigma\sigma$ be a name.

- a) Find a term $\text{injective} : (\sigma\sigma)B$ such that a logical interpretation satisfies the formula $\text{injective } f$ if and only if it interprets f as an injective function.
- b) Find a term $\text{surjective} : (\sigma\sigma)B$ such that a logical interpretation satisfies the formula $\text{surjective } f$ if and only if it interprets f as a surjective function.
- c) Find a formula finite that is satisfied by a logical interpretation \mathcal{I} if and only if $\mathcal{I}\sigma$ is a finite set.

Exercise 3.10 (Lists) Let the names $\text{nil} : L$, $\text{cons} : \sigma LL$, $\text{hd} : L\sigma$, and $\text{tl} : LL$ be given. Find a formula that is satisfied by a logical interpretation \mathcal{I} if and only if L represents all lists over σ and nil , cons , hd , and tl represent the list operations. Make sure that L contains no junk elements.