

**Lemma 2.2.1 (Generalized Cut)** If  $A \Rightarrow s$  and  $B \Rightarrow u$ , then  $A \# (B \setminus s) \Rightarrow u$ .

**Proof** If we try an induction on  $A \Rightarrow s$ , all cases go through but the case for the right implication rule. The problem can be fixed by adding an outermost induction on  $s$ , which gives us strong inductive hypotheses for the case  $s = (s_1 \rightarrow s_2)$ .

For each case of the inner inductions on  $A \Rightarrow s$  we have the assumption  $B \Rightarrow u$  and the claim  $A \# (B \setminus s) \Rightarrow u$ .

Let  $A \Rightarrow s$  be obtained with the variable rule. Then  $s = x \in A$ . Hence  $B \subseteq A \# (B \setminus s)$  and the claim follows by weakening from  $B \Rightarrow u$ .

Let  $A \Rightarrow s$  be obtained with the falsity rule. Then  $\perp \in A \# (B \setminus s)$  and the claim follows with the falsity rule.

Let  $A \Rightarrow s$  be obtained with the left implication rule. Then  $t_1 \rightarrow t_2 \in A$ ,  $A \Rightarrow t_1$ , and  $A, t_2 \Rightarrow s$ . The inductive hypothesis for  $A, t_2 \Rightarrow s$  is  $A, t_2 \# (B \setminus s) \Rightarrow u$ . We obtain the claim with the left implication rule for  $t_1 \rightarrow t_2$ , which leaves us with the proof obligations  $A \# (B \setminus s) \Rightarrow t_1$  and  $A \# (B \setminus s), t_2 \Rightarrow u$ . The obligations are easily obtained with weakening.

Let  $A \Rightarrow s$  be obtained with the right implication rule. We have  $s = (s_1 \rightarrow s_2)$  and  $A, s_1 \Rightarrow s_2$ . The inductive hypothesis for  $A, s_1 \Rightarrow s_2$  will not be used. The inductive hypotheses for  $s_1$  and  $s_2$  are as follows:

- For all  $A, B$  and  $u$ , if  $A \Rightarrow s_1$  and  $B \Rightarrow u$ , then  $A \# B \setminus s_1 \Rightarrow u$ .
- For all  $A, B$  and  $u$ , if  $A \Rightarrow s_2$  and  $B \Rightarrow u$ , then  $A \# B \setminus s_2 \Rightarrow u$ .

We prove the claim  $A \# B \setminus s \Rightarrow u$  by induction on  $B \Rightarrow u$ . This is the final induction.

We consider the case for the left implication rule, the other cases are straightforward. We have  $t_1 \rightarrow t_2 \in B$ ,  $B \Rightarrow t_1$ , and  $B, t_2 \Rightarrow u$ . The inductive hypotheses for  $B \Rightarrow t_1$  and  $B, t_2 \Rightarrow u$  are as follows:

- IH1:  $A \# B \setminus s \Rightarrow t_1$
- IH2:  $A \# (B, t_2) \setminus s \Rightarrow u$

We prove the claim  $A \# B \setminus s \Rightarrow u$  by case analysis.

1. Let  $s_1 = t_1$  and  $s_2 = t_2$ . The application of the inductive hypothesis for  $s_1$  to IH1 and  $A, s_1 \Rightarrow s_2$  yields

$$A \# B \setminus s \# (A, s_1) \setminus s_1 \Rightarrow s_2$$

The application of the inductive hypothesis for  $s_2$  to the above statement and IH2 yields

$$A \# B \setminus s \# (A, s_1) \setminus s_1 \# (A \# (B, s_2) \setminus s) \setminus s_2 \Rightarrow u$$

The claim now follows with weakening.

2. Let  $s \neq (t_1 \rightarrow t_2)$ . Then  $(t_1 \rightarrow t_2) \in A \# B \setminus s$ . We obtain the claim with the left implication rule. This leaves us with the proof obligations  $A \# B \setminus s \Rightarrow t_1$  and  $A \# B \setminus s, t_2 \Rightarrow u$ , which follow with IH1, IH2, and weakening. ■