Copying Garbage Collection

Guido Tack 14. November 2001 tack@ps.uni-sb.de

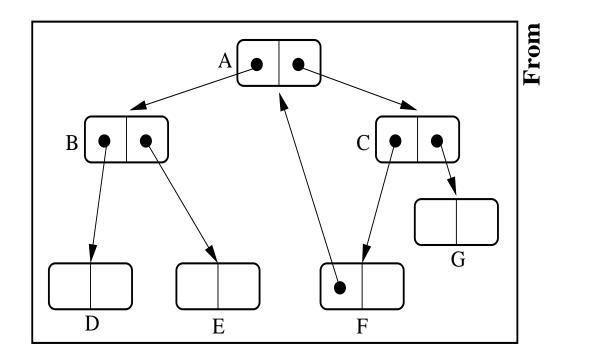
• Two "semi-spaces" (*From*-space and *To*-space)

- Two "semi-spaces" (*From*-space and *To*-space)
- Only From-space "active"

- Two "semi-spaces" (*From*-space and *To*-space)
- Only From-space "active"
- At GC time, copy the live nodes from From-Space to To-Space

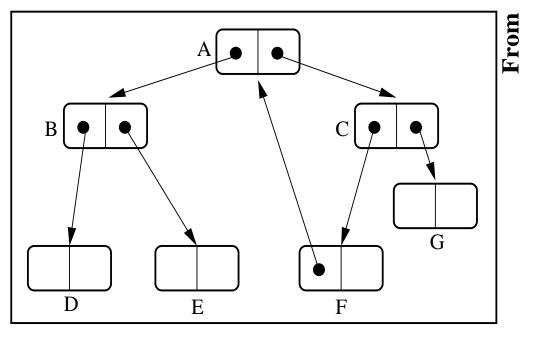
- Two "semi-spaces" (*From*-space and *To*-space)
- Only From-space "active"
- At GC time, copy the live nodes from From-Space to To-Space
- Then "flip" the spaces

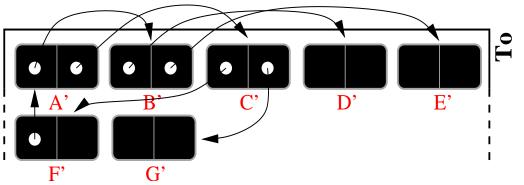
2 Before GC





3 AFTER GC





• Iterative algorithm

- Iterative algorithm
- Interleaves copying and scanning

- Iterative algorithm
- Interleaves copying and scanning
- Two pointers needed: scan / free

- Iterative algorithm
- Interleaves copying and scanning
- Two pointers needed: scan / free
- Forwarding pointers used to preserve sharing

Tricolour abstraction:

Tricolour abstraction:

• Black nodes: GC finished, not to be considered again

Tricolour abstraction:

- Black nodes: GC finished, not to be considered again
- Grey nodes: Visited but not completed

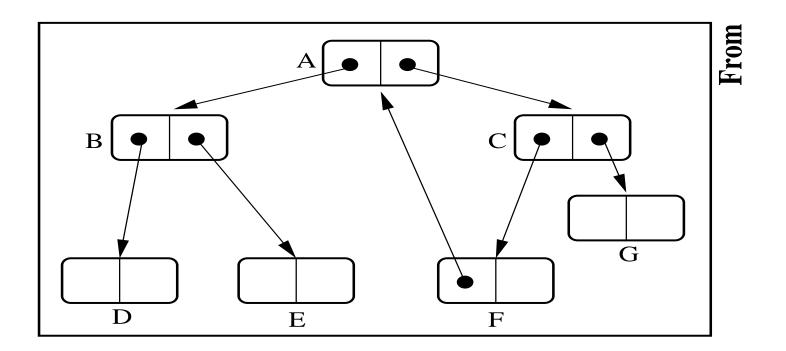
Tricolour abstraction:

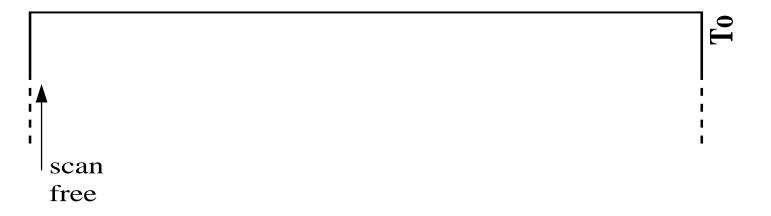
- Black nodes: GC finished, not to be considered again
- Grey nodes: Visited but not completed
- White nodes: Unvisited, considered garbage after tracing

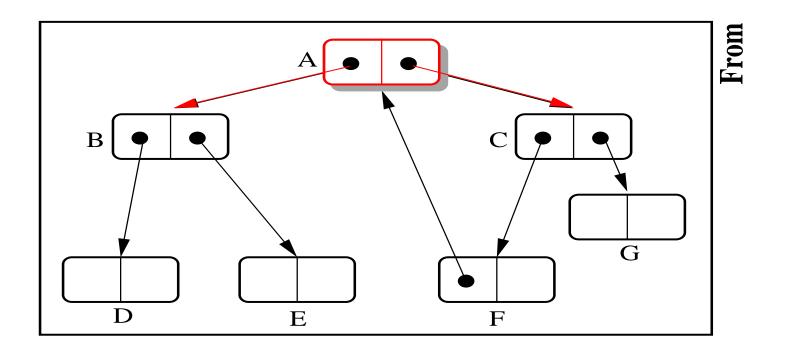
Tricolour abstraction:

- Black nodes: GC finished, not to be considered again
- Grey nodes: Visited but not completed
- White nodes: Unvisited, considered garbage after tracing

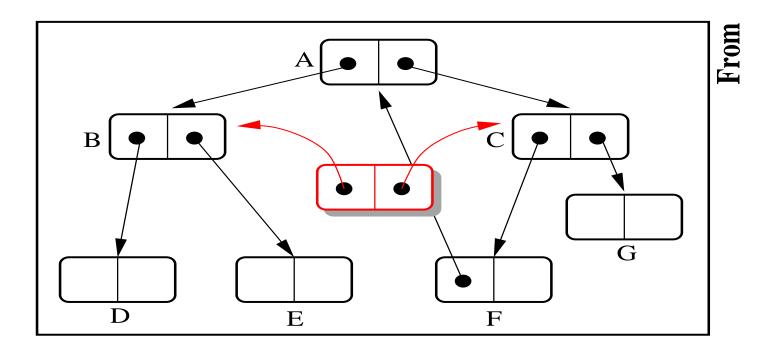
GC terminates when all reachable nodes are black



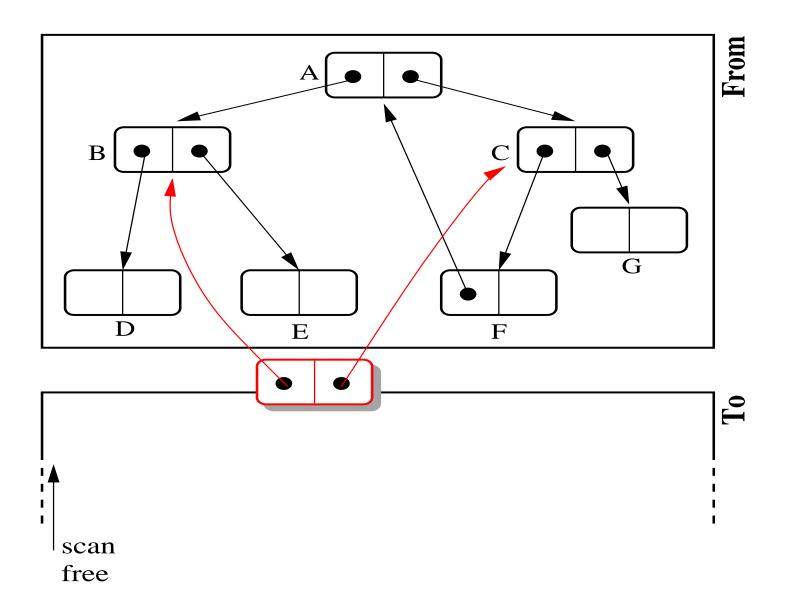


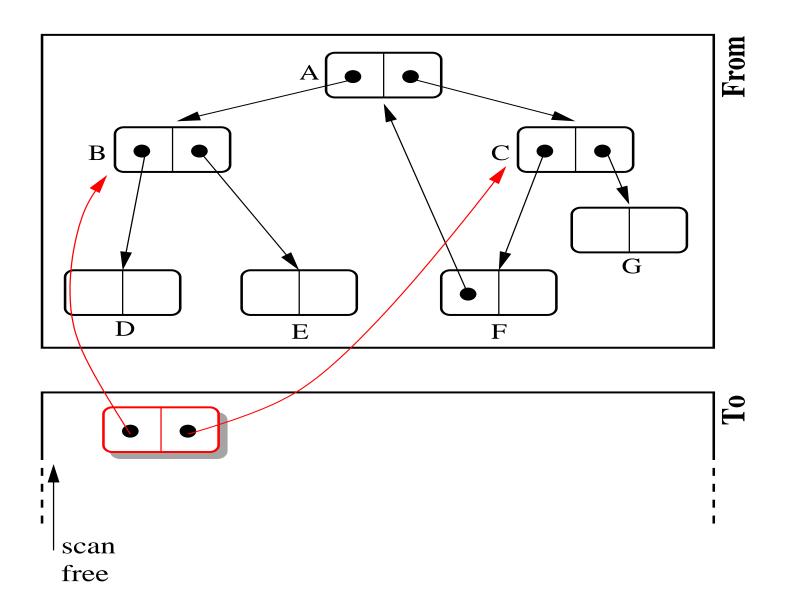


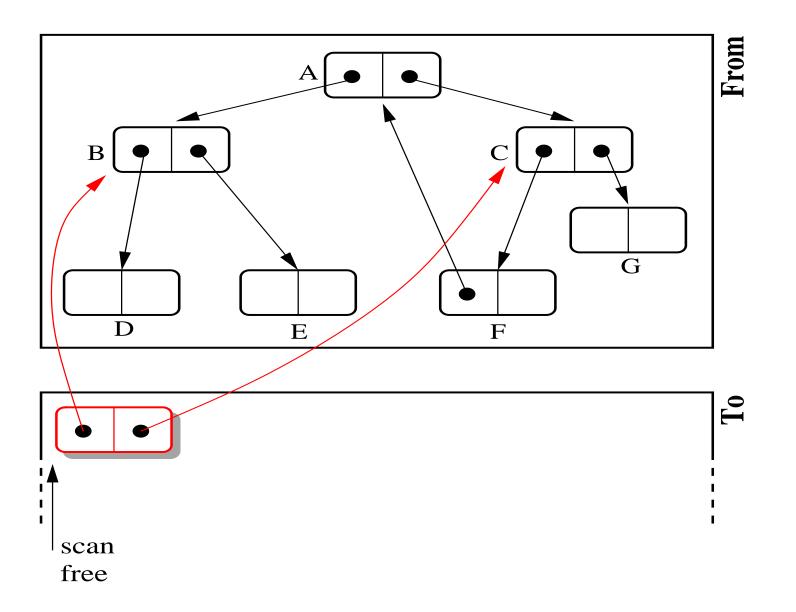


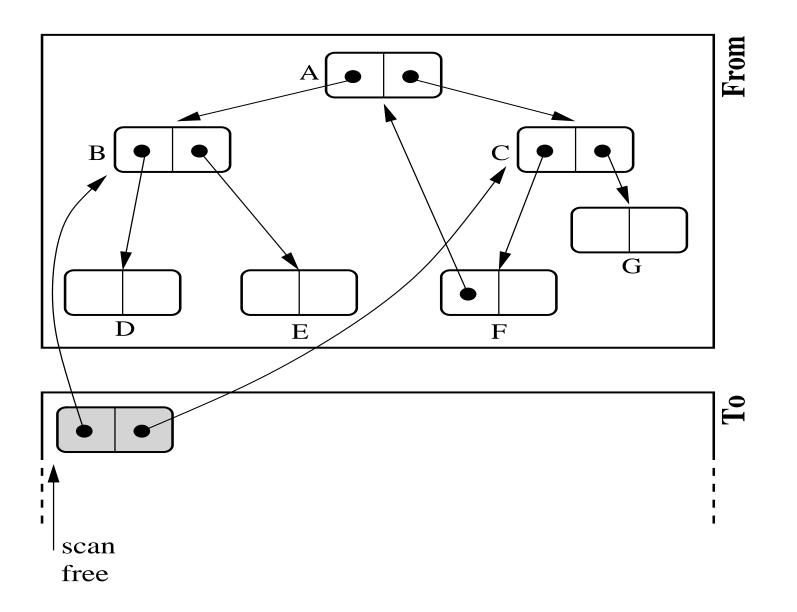


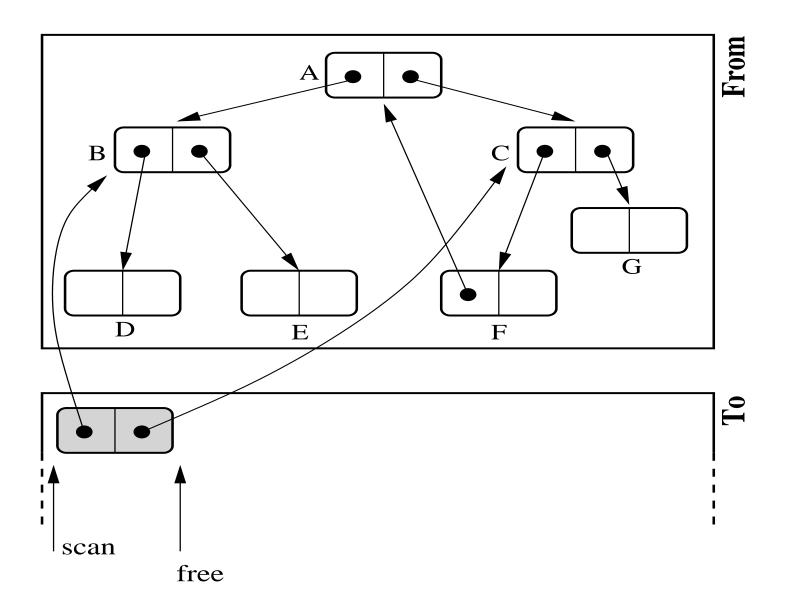


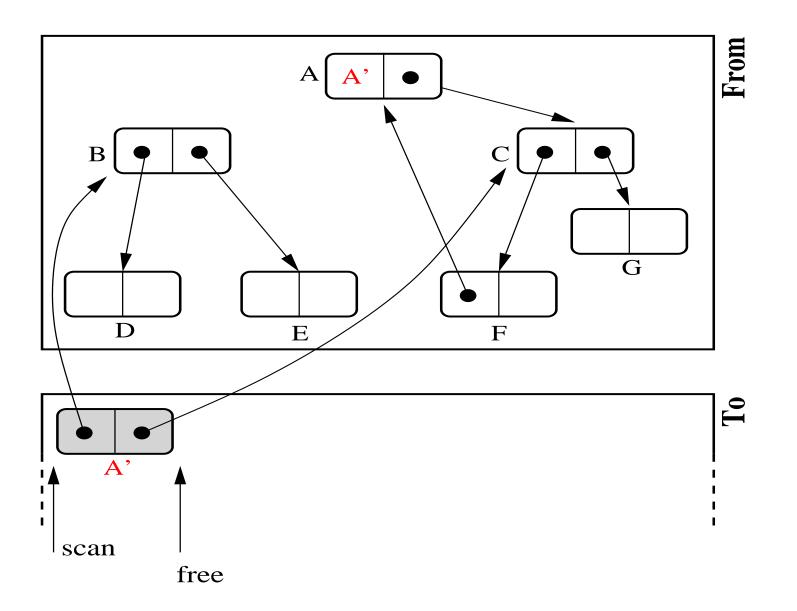


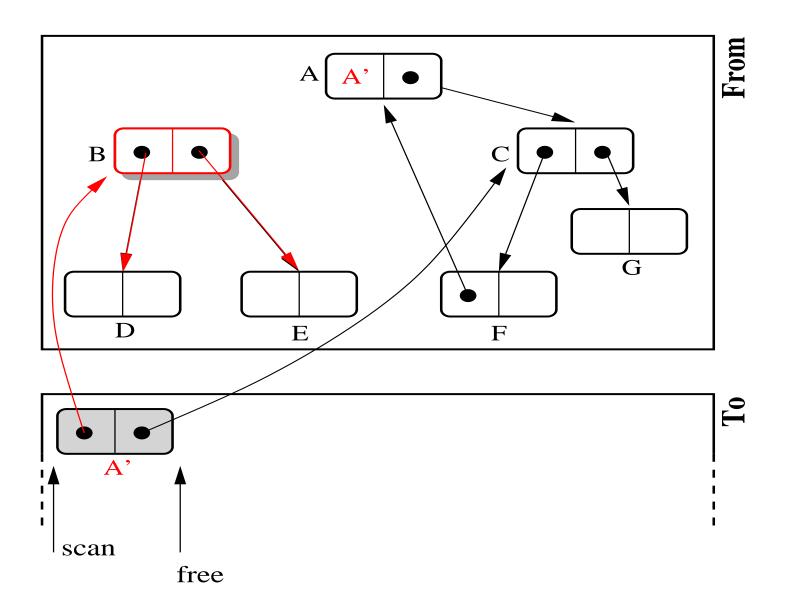


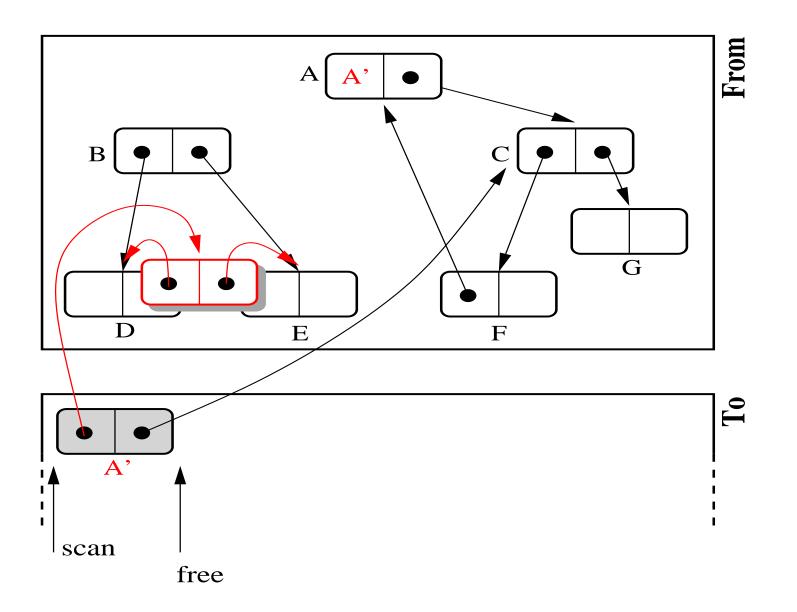


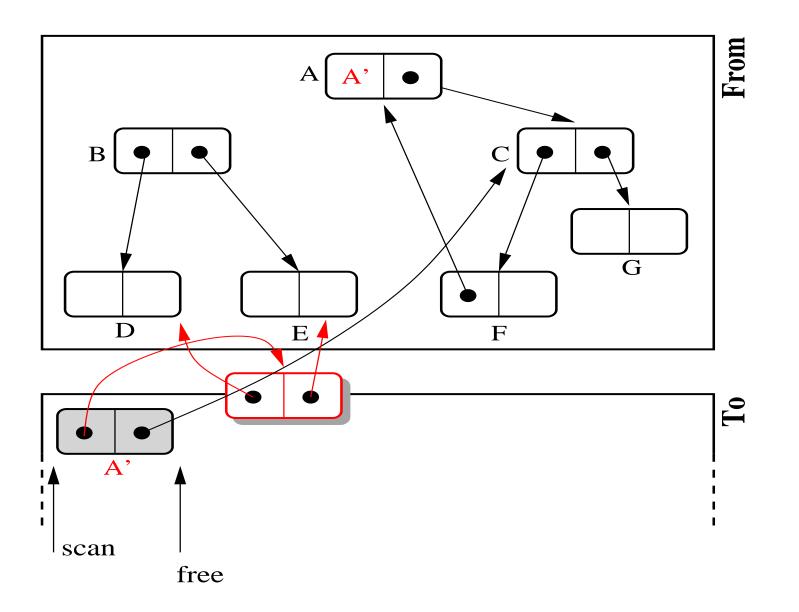


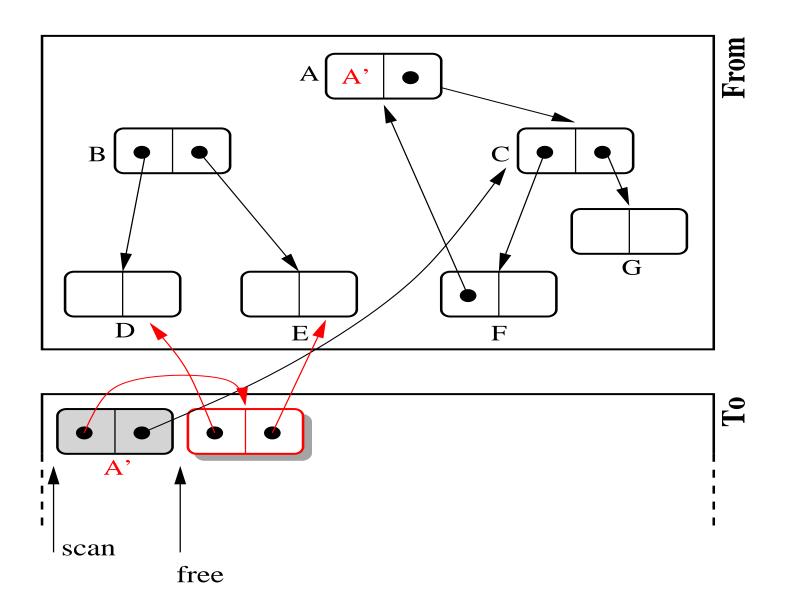


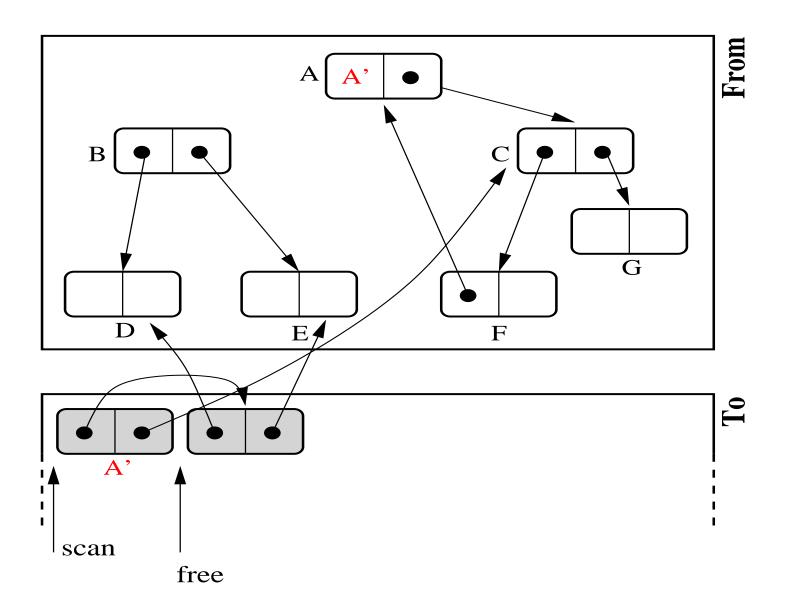


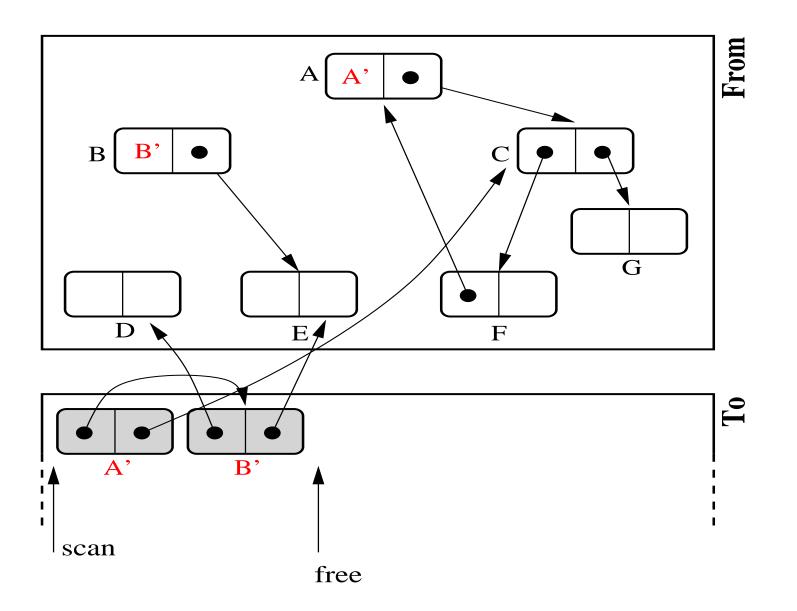


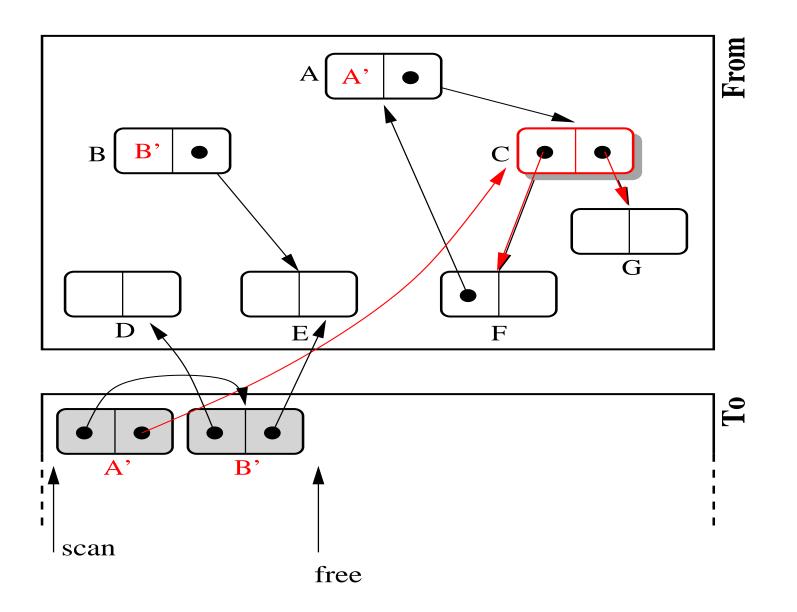


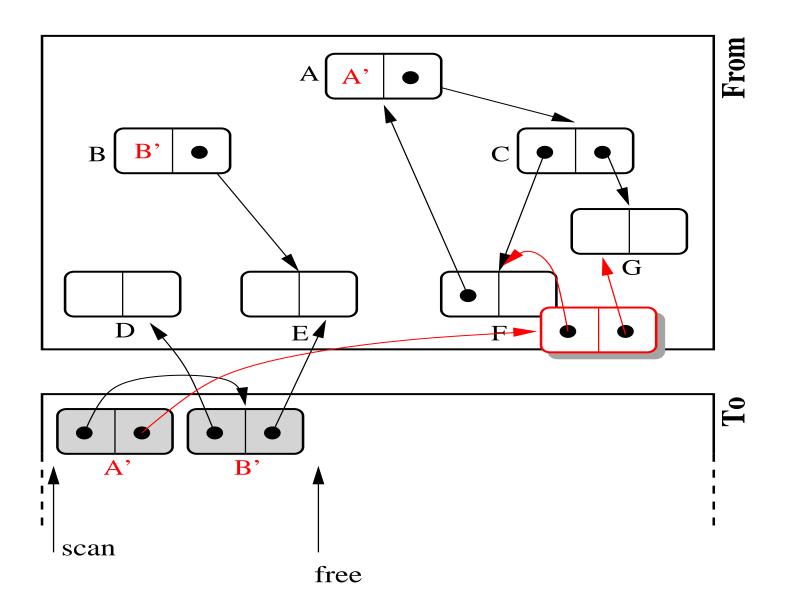


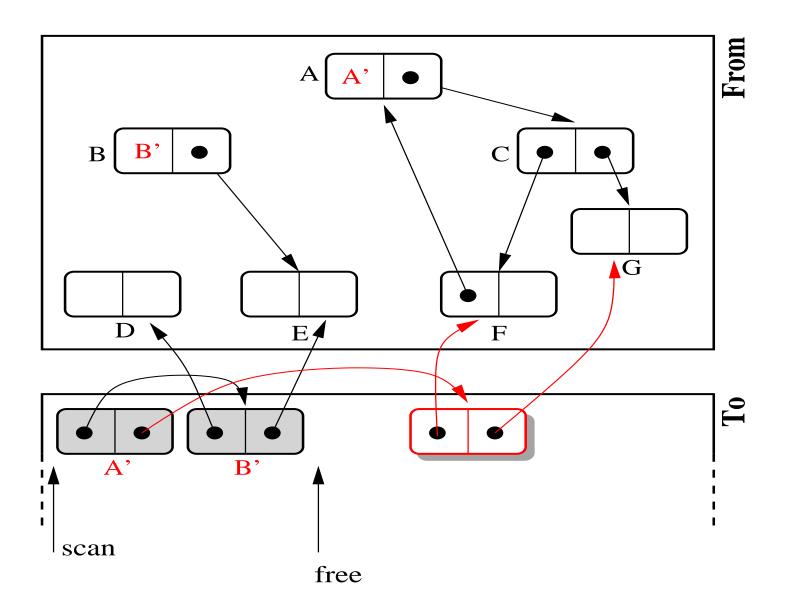


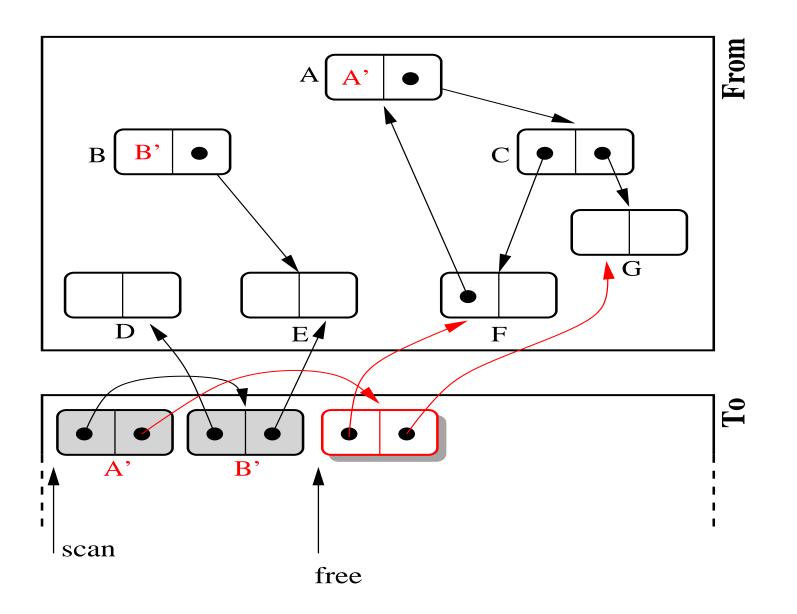


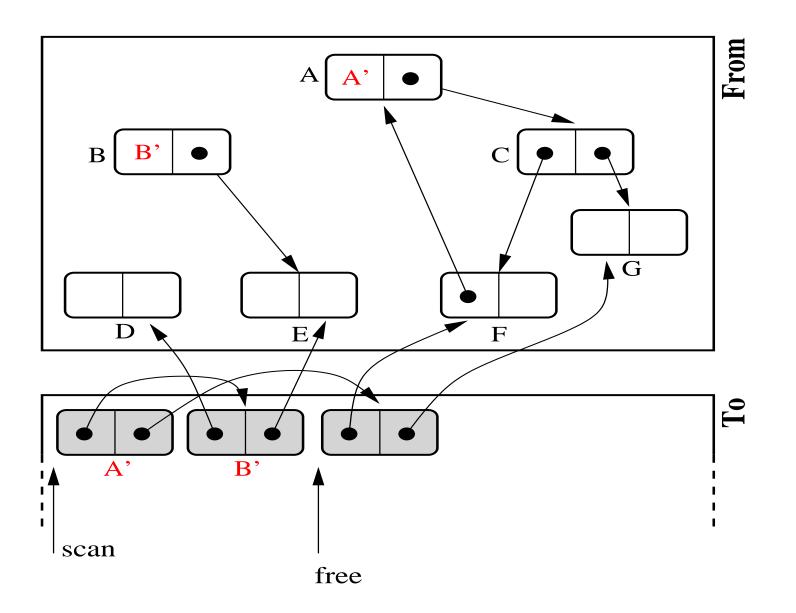


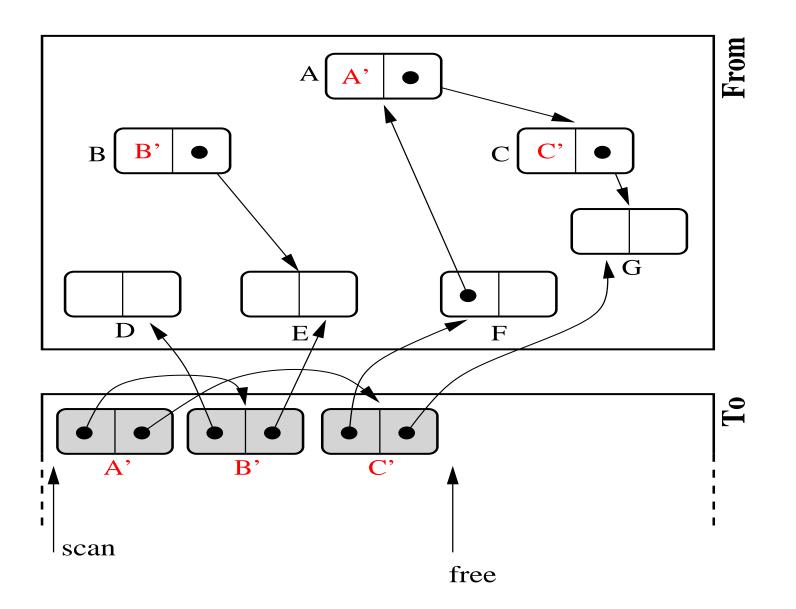


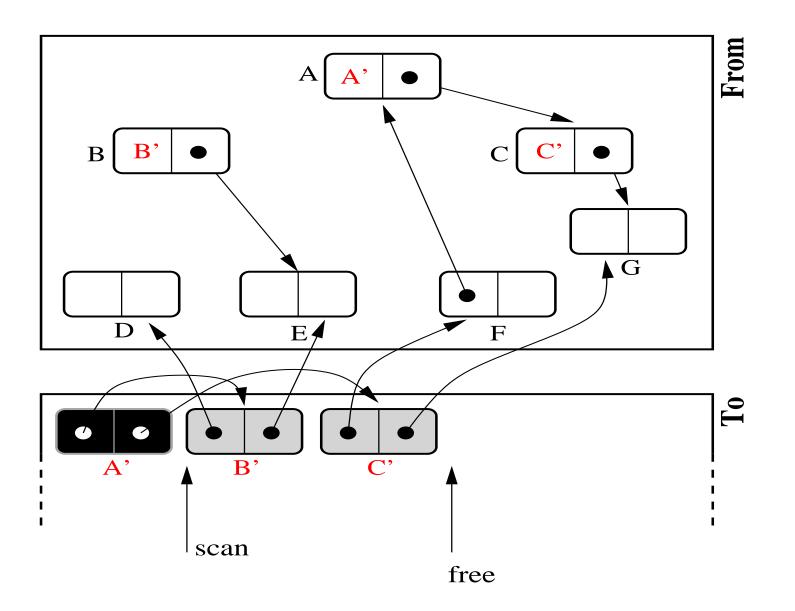


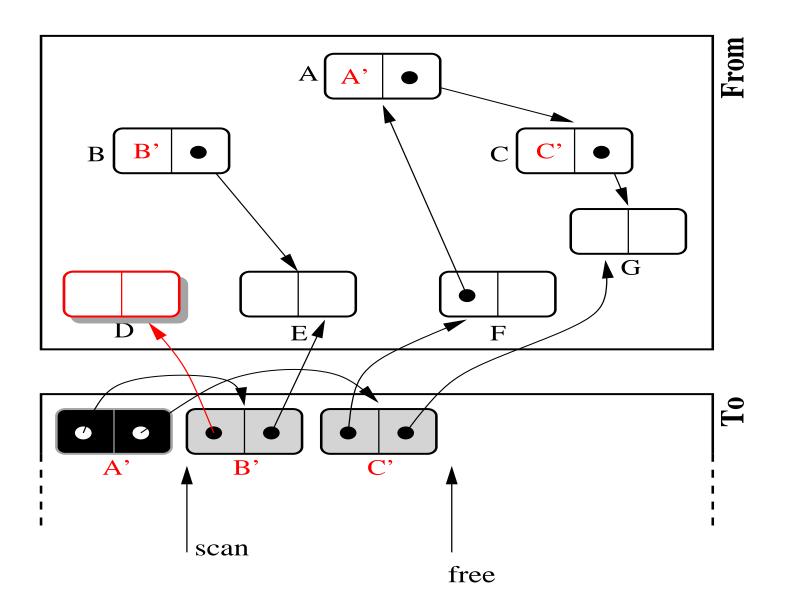


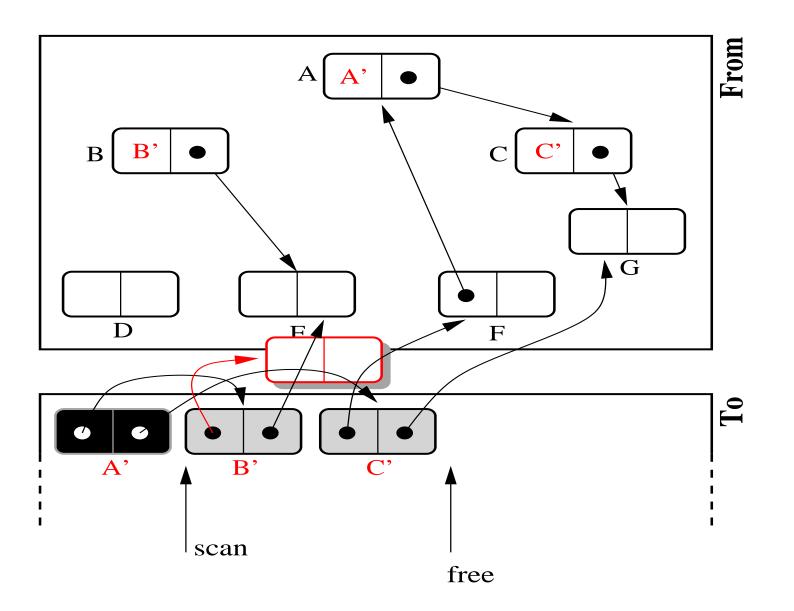


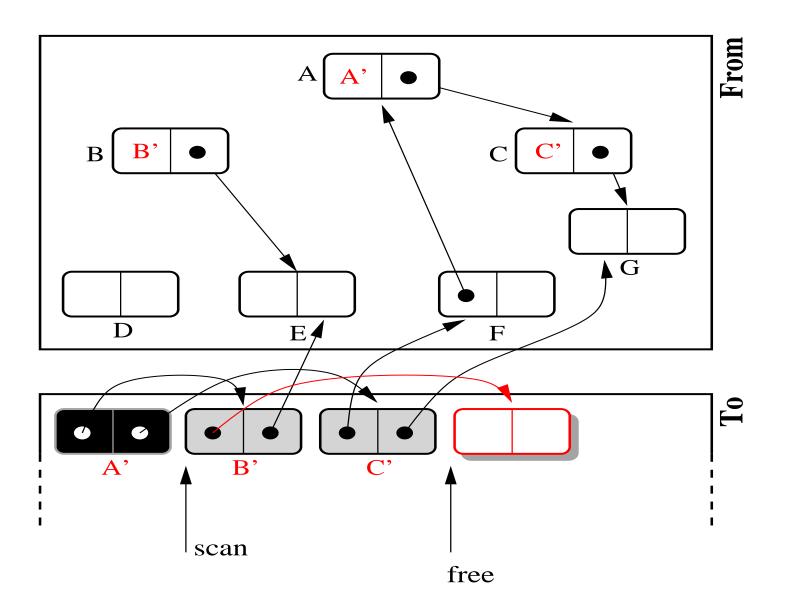


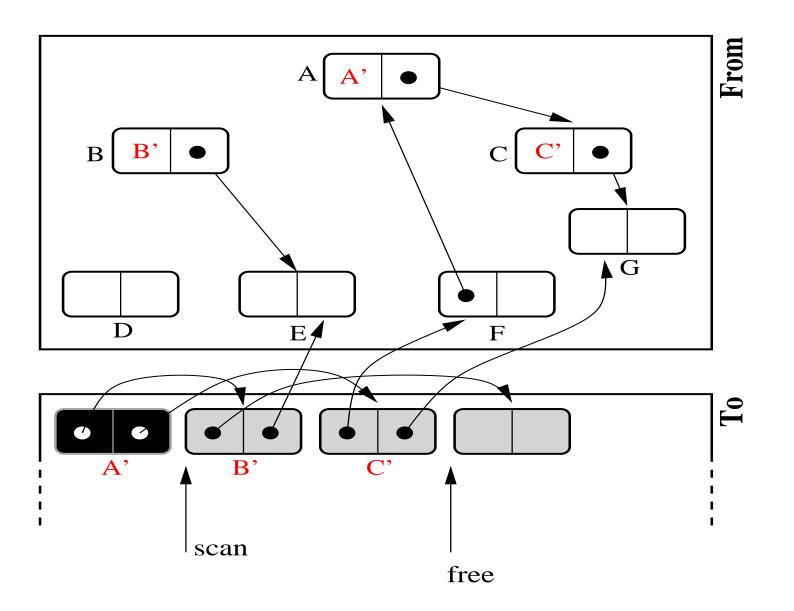


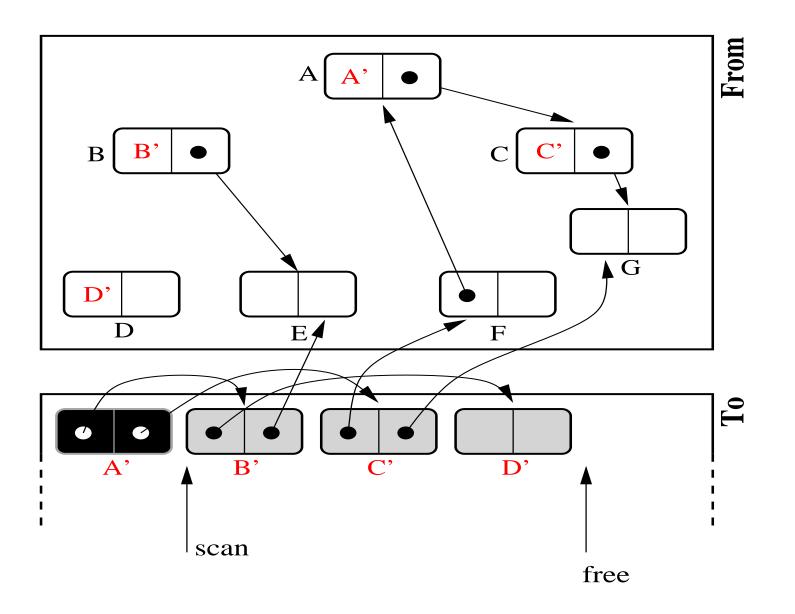


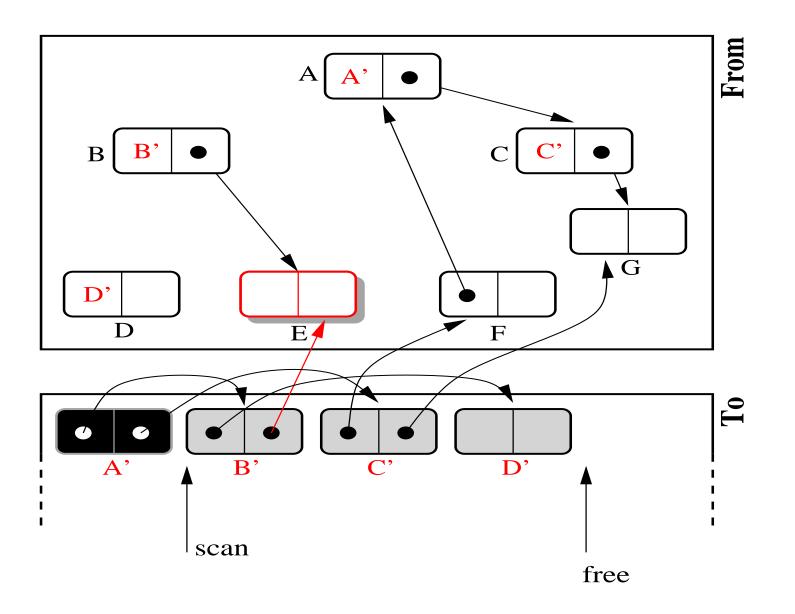


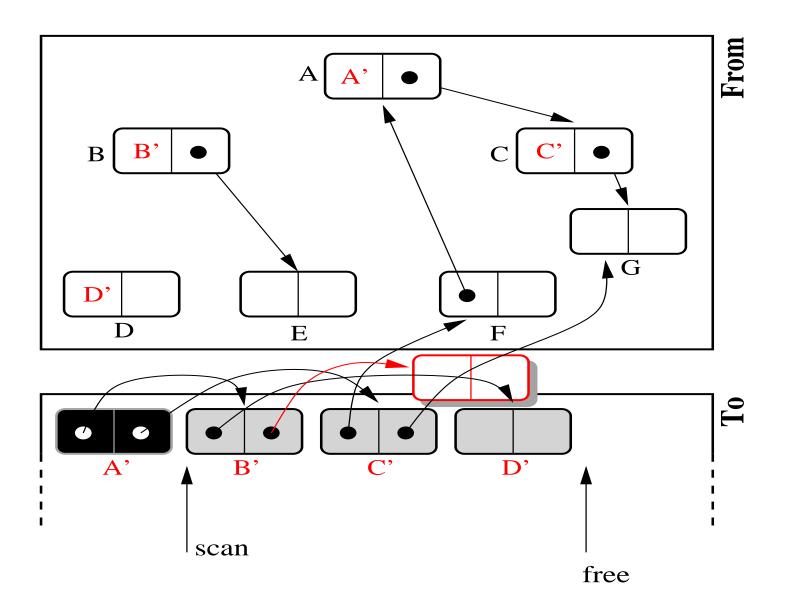


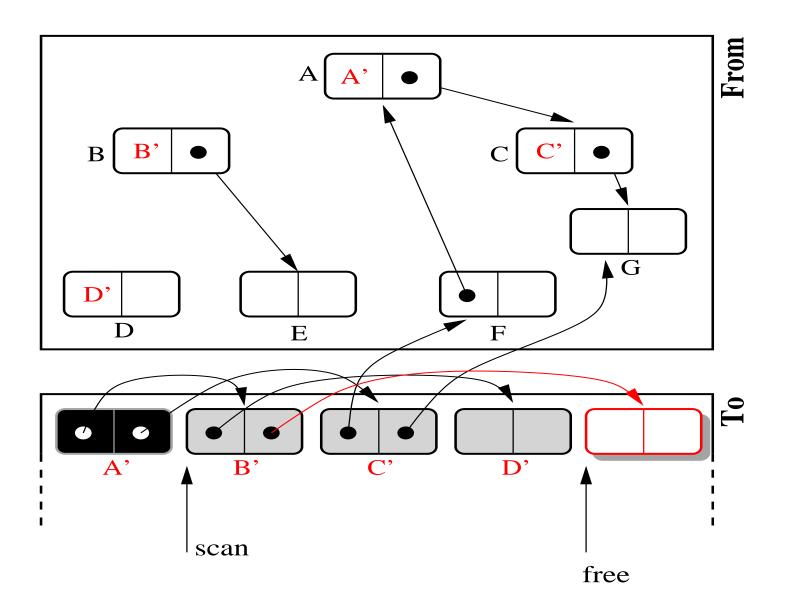


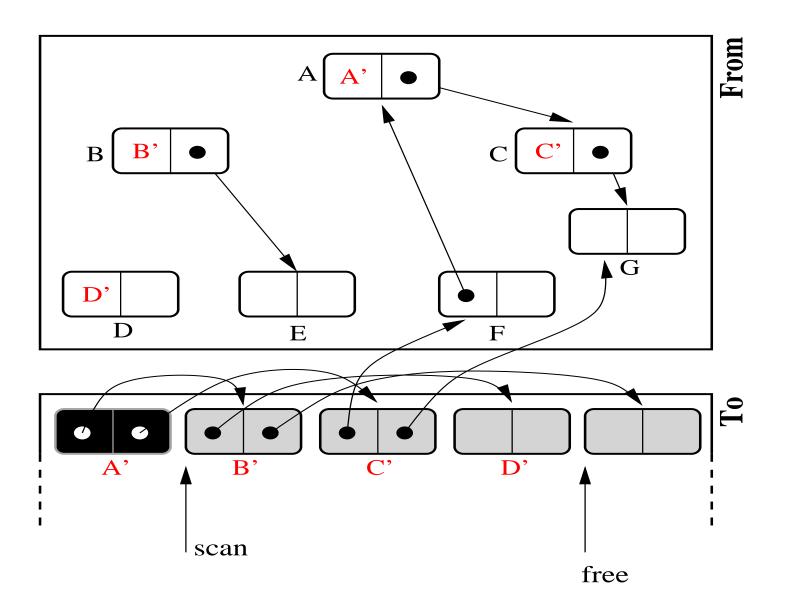


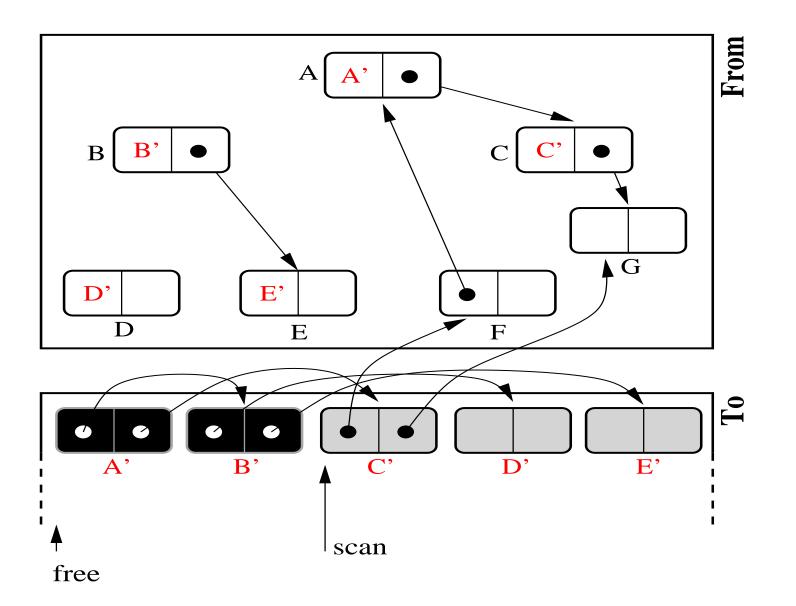


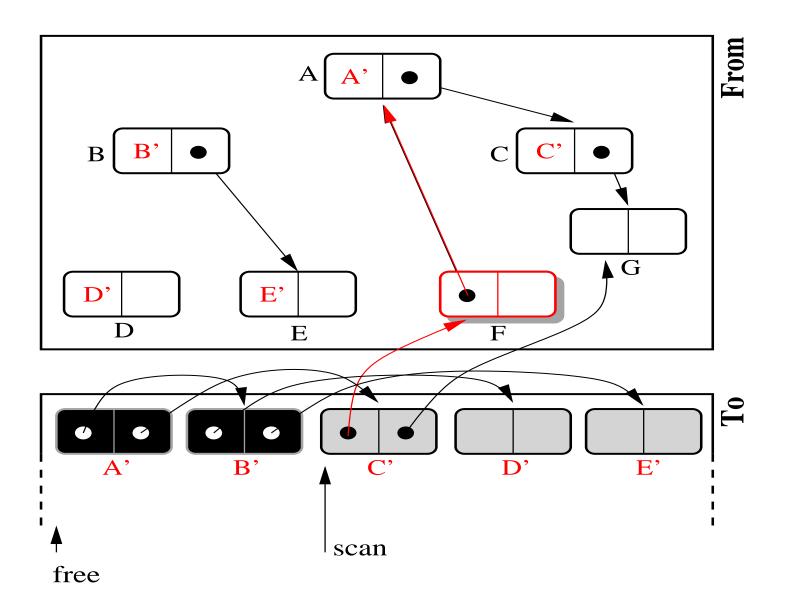


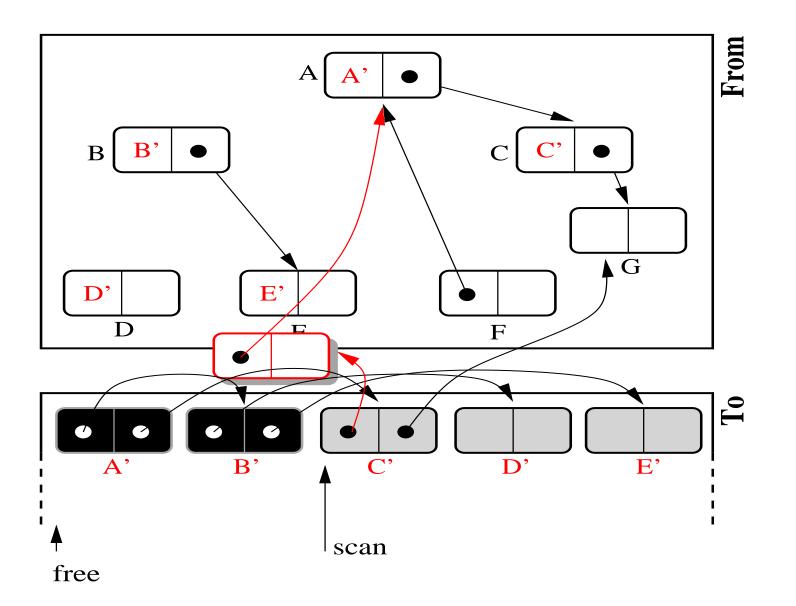


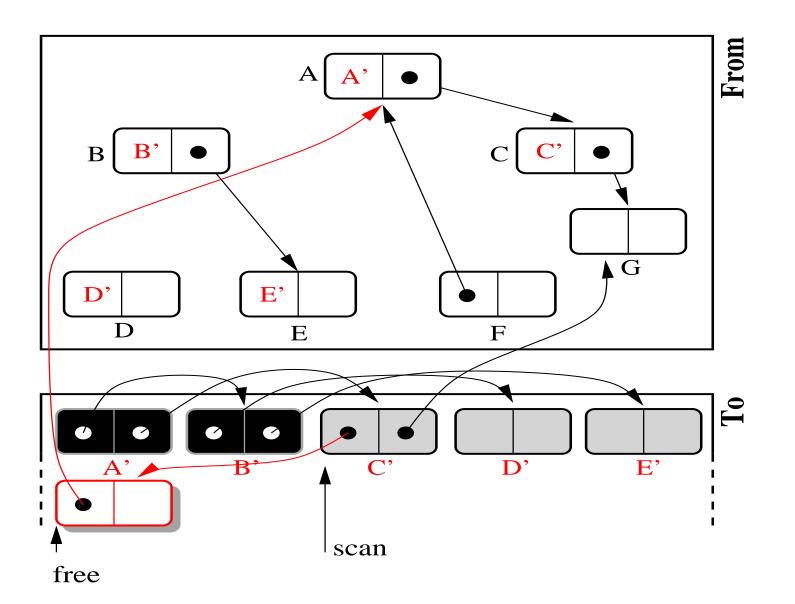


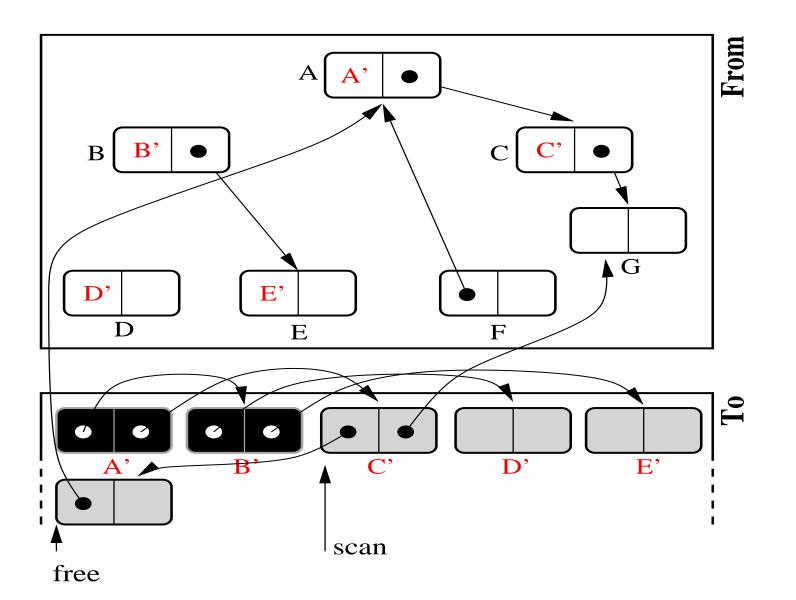


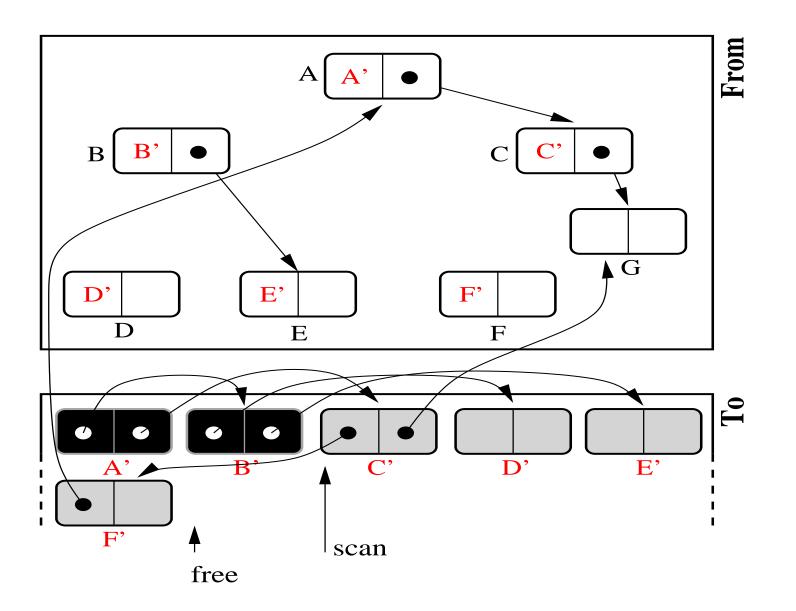


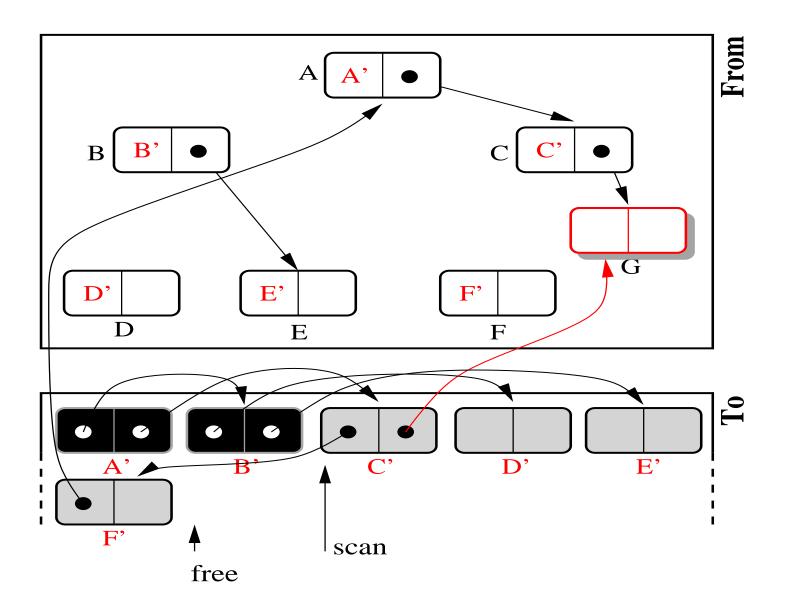


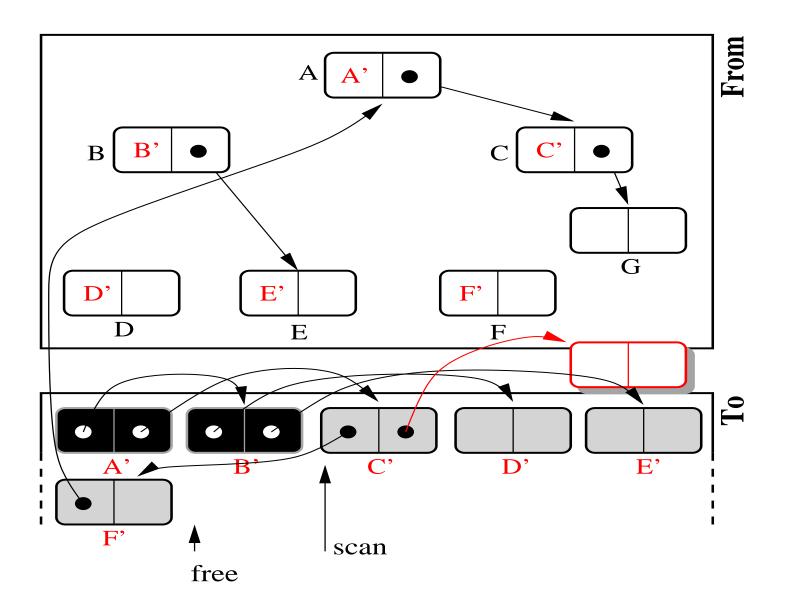


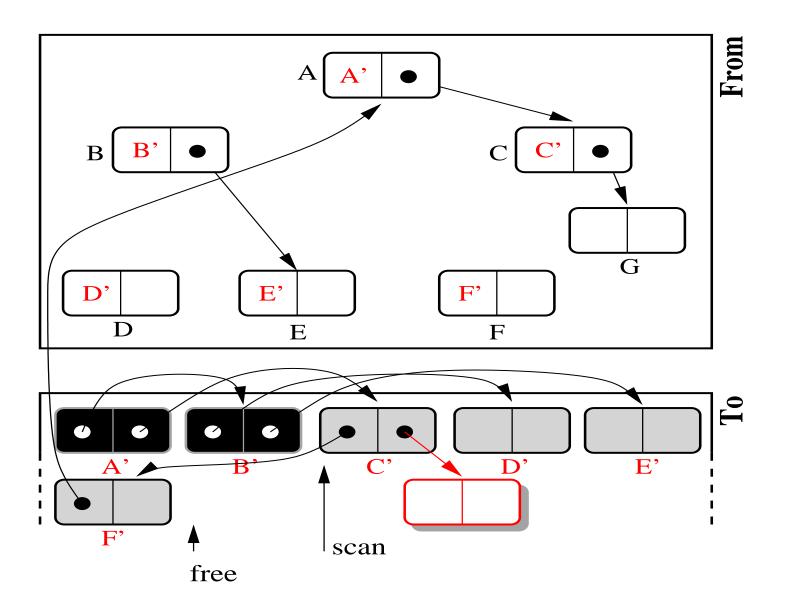


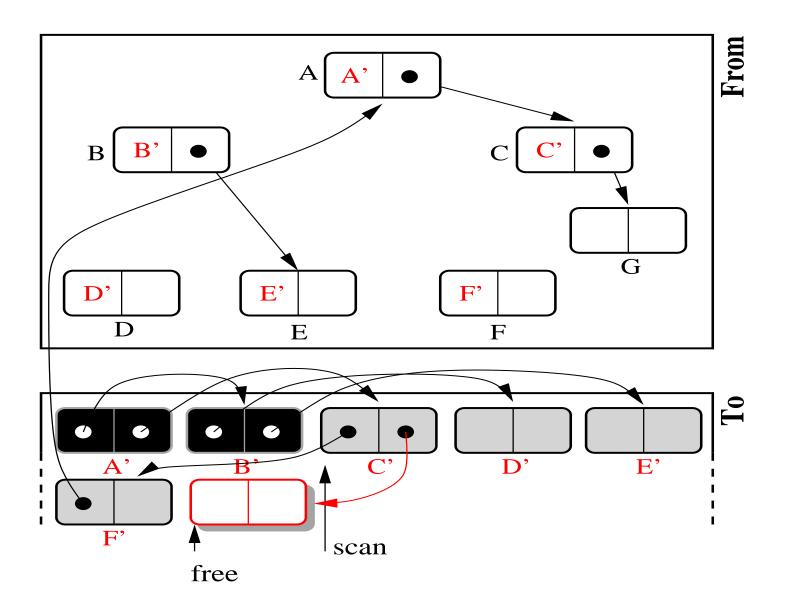


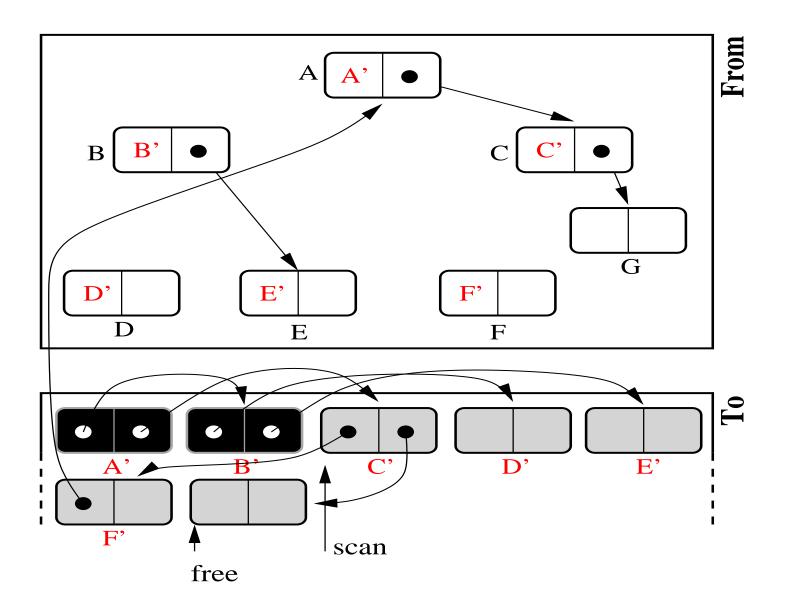


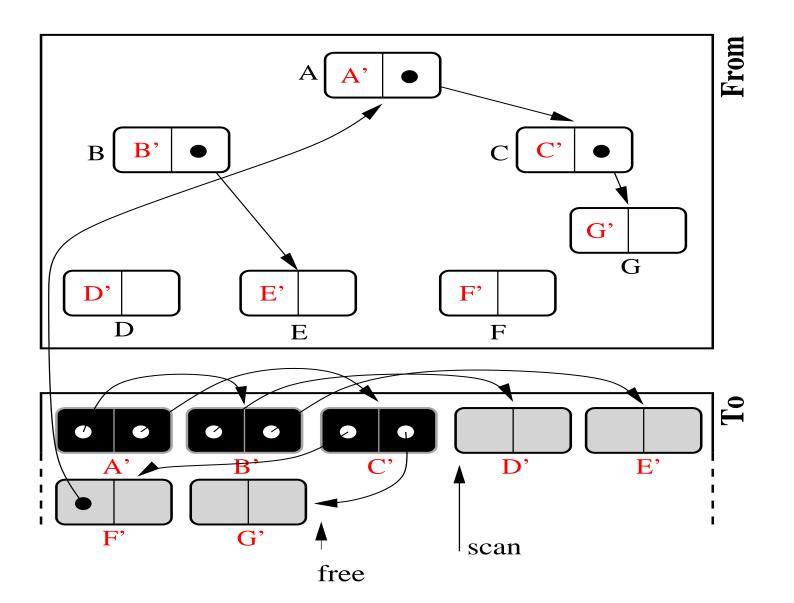


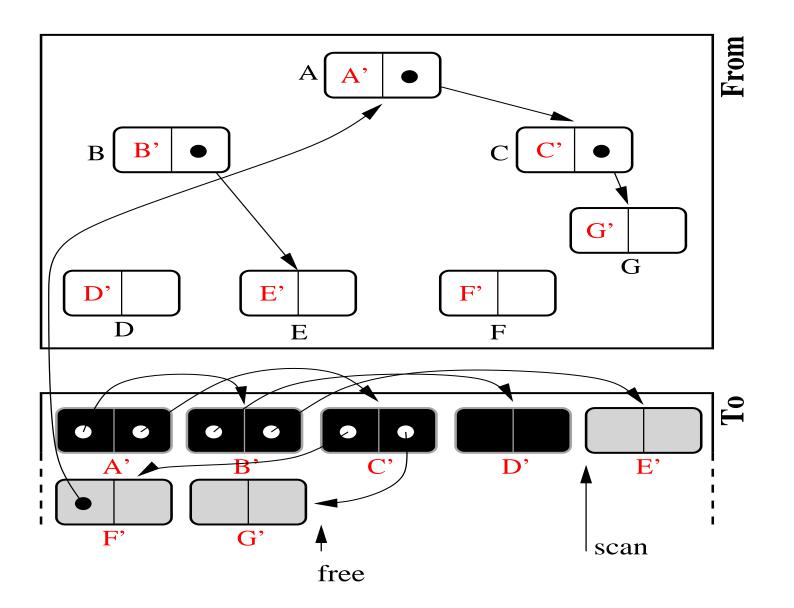


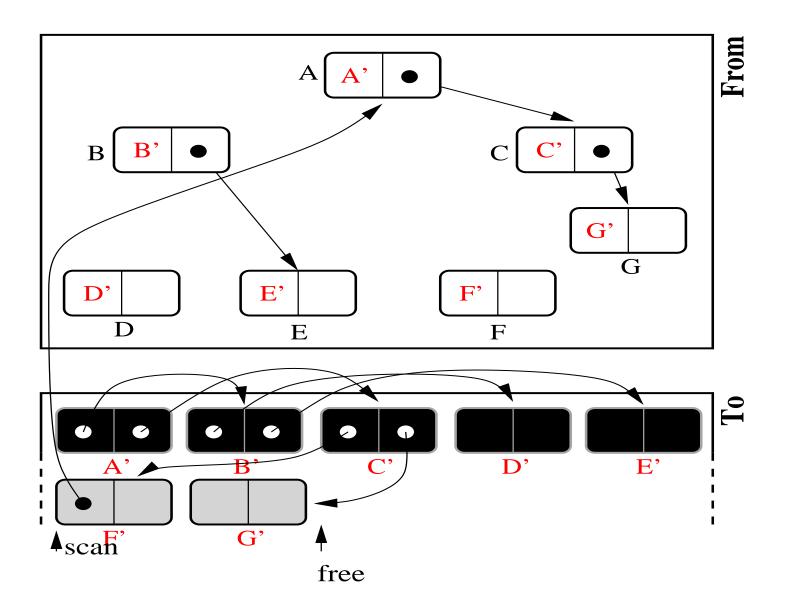


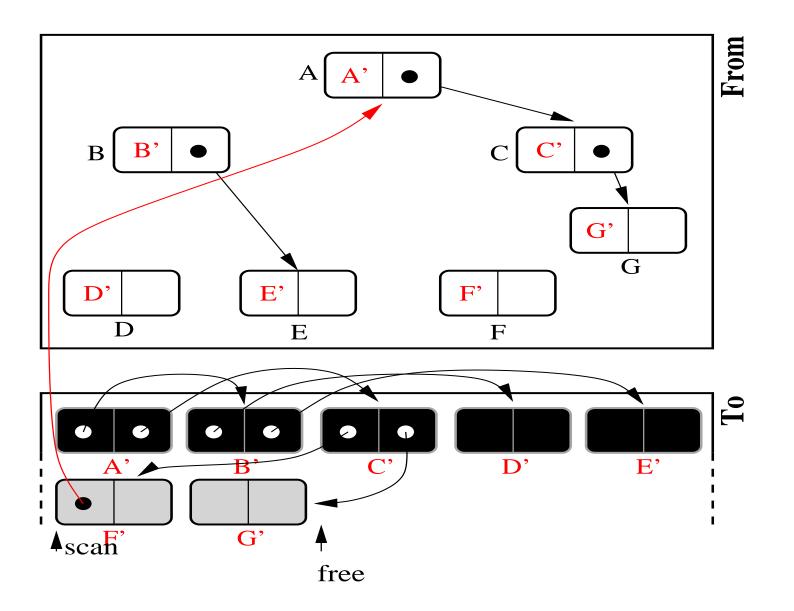


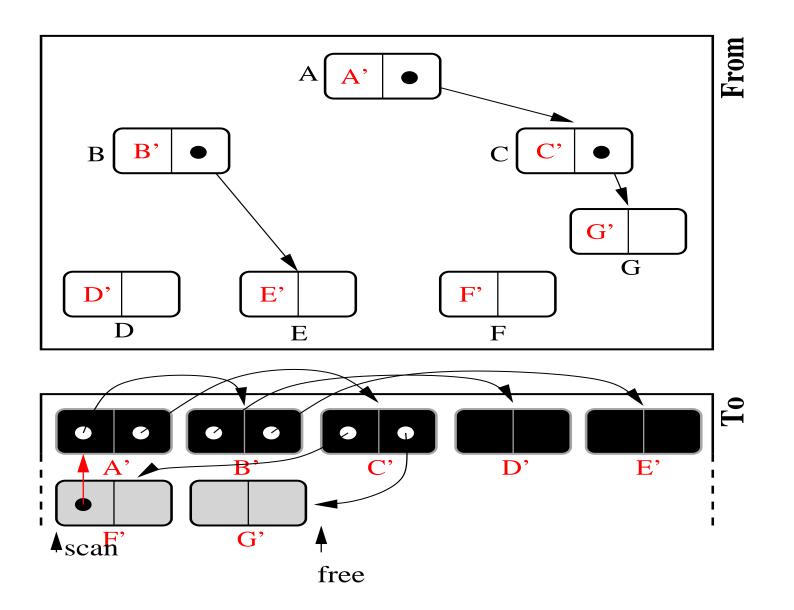


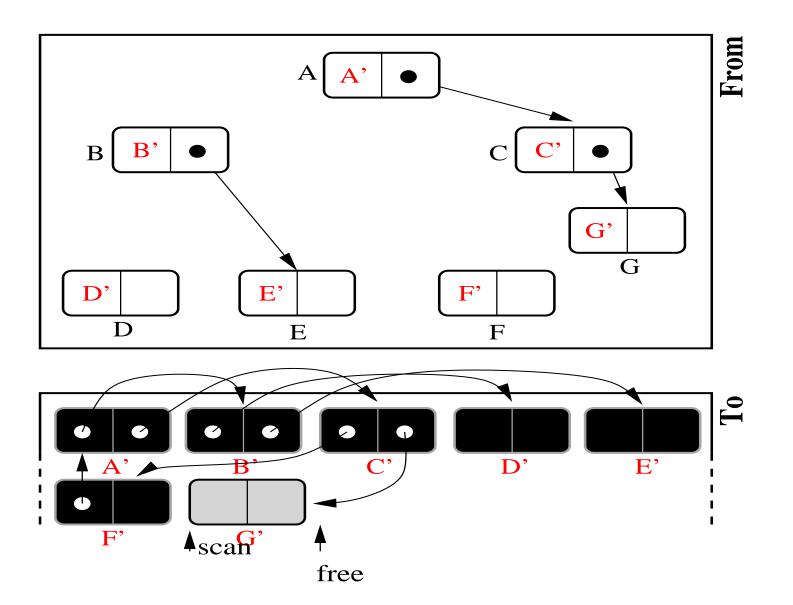


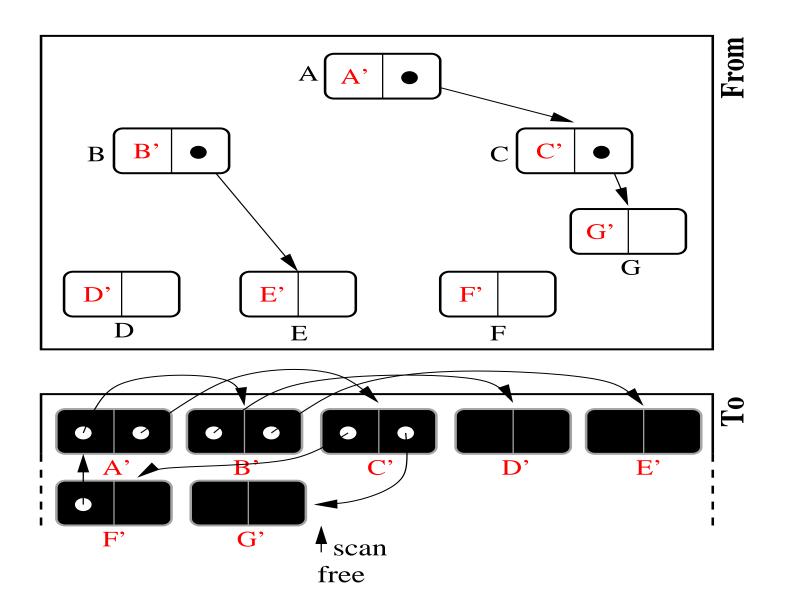












• Size of nodes known

- Size of nodes known
- Child fields known

- Size of nodes known
- Child fields known
- Two logically contiguous areas for the heap

6 Assumptions

- Size of nodes known
- Child fields known
- Two logically contiguous areas for the heap
- Ability to mark nodes

• Constant stack required

7 PROPERTIES

- Constant stack required
- Time depends only on no. of *live* nodes

- Constant stack required
- Time depends only on no. of *live* nodes
- Covers cycles / sharing

- Constant stack required
- Time depends only on no. of *live* nodes
- Covers cycles / sharing
- Heap is compacted

- Constant stack required
- Time depends only on no. of *live* nodes
- Covers cycles / sharing
- Heap is compacted
- Stop/start collector

- Constant stack required
- Time depends only on no. of *live* nodes
- Covers cycles / sharing
- Heap is compacted
- Stop/start collector
- Nothing to do between GCs

High Level

High Level

• Copying large objects is expensive

High Level

- Copying large objects is expensive
- Twice the logical memory needed

High Level

- Copying large objects is expensive
- Twice the logical memory needed

Low Level

• Breadth first \Rightarrow decreased locality

High Level

- Copying large objects is expensive
- Twice the logical memory needed

- Breadth first \Rightarrow decreased locality
- Paging issues

• Large Object Areas

• Large Object Areas (possibly handled by different collector)

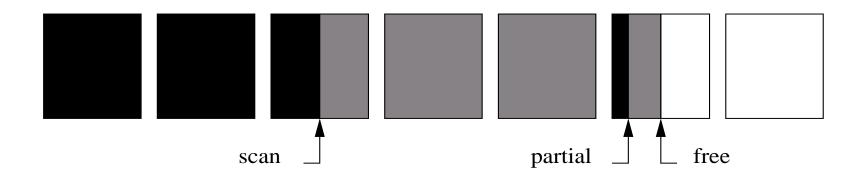
- Large Object Areas (possibly handled by different collector)
- Areas for long living objects

- Large Object Areas (possibly handled by different collector)
- Areas for long living objects (only scanned, not copied)

- Large Object Areas (possibly handled by different collector)
- Areas for long living objects (only scanned, not copied)
- Maintain locality by using other exploration strategies

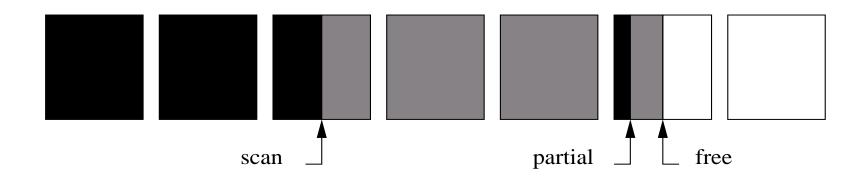
- Large Object Areas (possibly handled by different collector)
- Areas for long living objects (only scanned, not copied)
- Maintain locality by using other exploration strategies (but then stack becomes an issue)

Approximately depth-first copying(1)



Approximately depth-first copying(1)

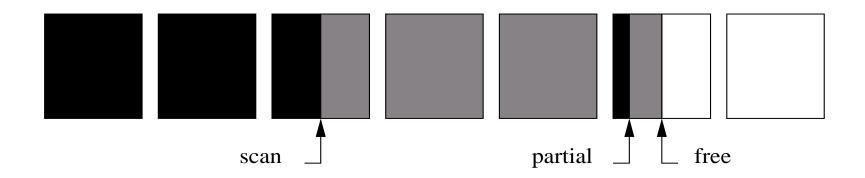
Modification of Cheney's agorithm by Moon (1984)



Approximately depth-first copying(1)

Modification of Cheney's agorithm by Moon (1984)

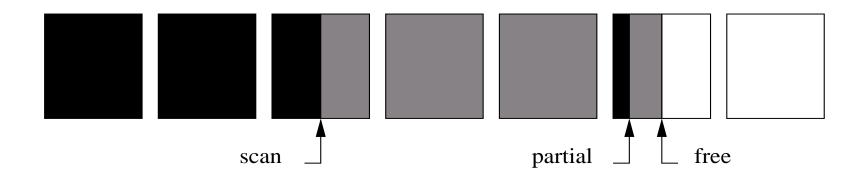
• Always start scanning on the last partially filled page in To-space



Approximately depth-first copying(1)

Modification of Cheney's agorithm by Moon (1984)

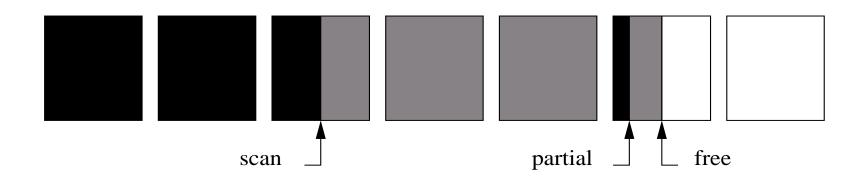
- Always start scanning on the last partially filled page in To-space
- When that page is completed, continue with ordinary scan



Approximately depth-first copying(1)

Modification of Cheney's agorithm by Moon (1984)

- Always start scanning on the last partially filled page in To-space
- When that page is completed, continue with ordinary scan
- As soon as an object is copied, start partial scan again



Approximately depth-first copying(2)

Approximately depth-first copying(2)

• Drawback: Some nodes are scanned twice

Approximately depth-first copying(2)

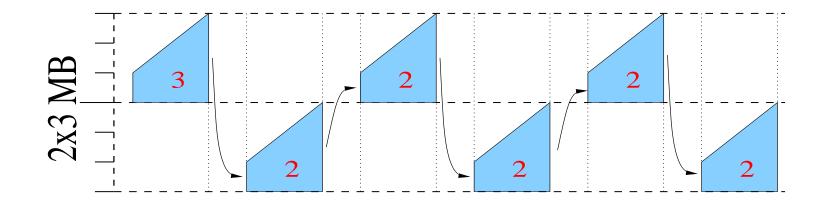
- Drawback: Some nodes are scanned twice
- Tests indicate a 15 percent improvement of locality

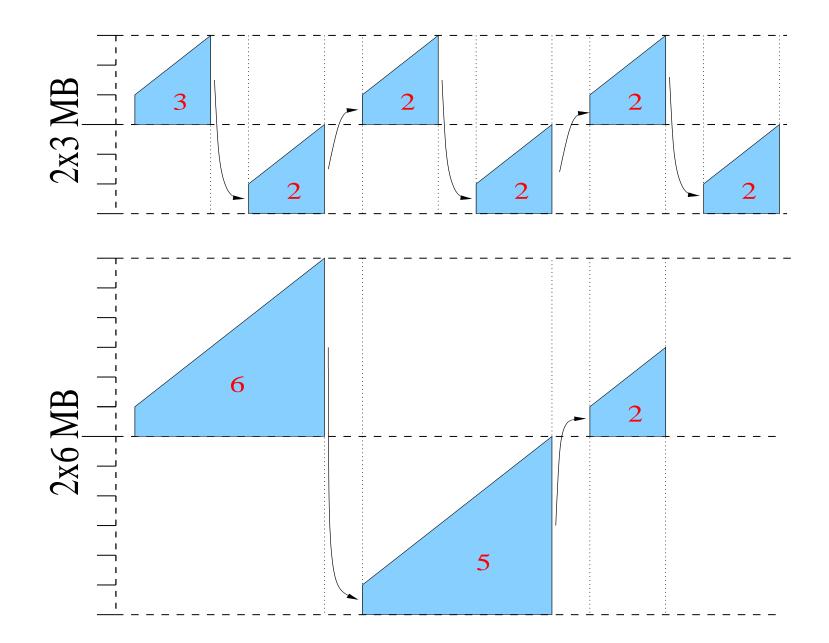
Approximately depth-first copying(2)

- Drawback: Some nodes are scanned twice
- Tests indicate a 15 percent improvement of locality
- ... and a 6 percent increased GC time

• Making the semi-spaces bigger decreases frequency of GC

- Making the semi-spaces bigger decreases frequency of GC
- Less frequent GC means older objects
 - $\Rightarrow \mathsf{More} \ \mathsf{garbage}$





• Storage management dominated by allocation (alloc. is cheap)

- Storage management dominated by allocation (alloc. is cheap)
- Many small, short-lived objects (copying small objects not much more expensive than marking)

- Storage management dominated by allocation (alloc. is cheap)
- Many small, short-lived objects (copying small objects not much more expensive than marking)
- GC delay doesn't matter (no real-time system)

• Hybrid systems can be used

 Hybrid systems can be used (e.g. use copying only for small objects, mark-sweep for large objects)

- Hybrid systems can be used (e.g. use copying only for small objects, mark-sweep for large objects)
- Copying collection can be used as a foundation for

GC algorithms.

- Hybrid systems can be used (e.g. use copying only for small objects, mark-sweep for large objects)
- Copying collection can be used as a foundation for
 - incremental
 - GC algorithms.

- Hybrid systems can be used (e.g. use copying only for small objects, mark-sweep for large objects)
- Copying collection can be used as a foundation for
 - incremental
 - generational
 - GC algorithms.