



# Semantics of Programming Languages Solutions to Assignment 8

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**Errata.** The projections  $\Pi_i$  should be functions from  $D = 2^{D_1 \times \dots \times D_n}$  to  $2^{D_i}$ , otherwise their definitions would be nonsense.

**Exercise 8.1** To prove that in general there is no such function  $f$  it suffices to show that there is none for  $n = 2$  and  $D_1 = D_2 = \{0, 1\}$ .

Let  $S = \{\langle 0, 0 \rangle, \langle 1, 1 \rangle\}$ . Then  $\text{app}(S) = \{0, 1\} \times \{0, 1\}$  is a set of cardinality 4 whereas for every function  $f : D_1 \times D_2 \rightarrow D_1 \times D_2$ , the set  $\{f(s) \mid s \in S\} = \{f(\langle 0, 0 \rangle), f(\langle 1, 1 \rangle)\}$  has cardinality 1 or 2.

## Exercise 8.2

1. Let  $S \in D$ . Then  $\text{app}(S) = \Pi_1(S) \times \dots \times \Pi_n(S) = \gamma(\langle \Pi_1(S), \dots, \Pi_n(S) \rangle) = \gamma(\alpha(S))$ .
2. Let  $S \in D$ . To prove that  $\alpha(S)$  is the least element in  $D^\#$  whose image under  $\gamma$  is above  $S$ , fix an arbitrary  $\langle M_1, \dots, M_n \rangle \in D^\#$  such that  $\gamma(\langle M_1, \dots, M_n \rangle) \supseteq S$ . Let  $x \in \Pi_i(S)$ , i. e., there exists  $\langle x_1, \dots, x_i, \dots, x_n \rangle \in S$  such that  $x = x_i$ . Then  $\langle x_1, \dots, x_i, \dots, x_n \rangle \in \gamma(\langle M_1, \dots, M_i, \dots, M_n \rangle) = M_1 \times \dots \times M_i \times \dots \times M_n$ , so  $x = x_i \in M_i$ . Hence  $\alpha(S) = \langle \Pi_1(S), \dots, \Pi_n(S) \rangle \sqsubseteq \langle M_1, \dots, M_n \rangle$ . Furthermore  $\gamma(\alpha(S)) \supseteq S$ , see exercise 8.3.2.
3. Let  $\langle M_1, \dots, M_n \rangle \in D^\#$ . To prove that  $\gamma(\langle M_1, \dots, M_n \rangle)$  is the greatest element in  $D$  whose image under  $\alpha$  is below  $\langle M_1, \dots, M_n \rangle$ , fix an arbitrary  $S \in D$  such that  $\alpha(S) \sqsubseteq \langle M_1, \dots, M_n \rangle$ . Since  $\langle \Pi_1(S), \dots, \Pi_n(S) \rangle = \alpha(S) \sqsubseteq \langle M_1, \dots, M_n \rangle$ , we have for all  $i$ ,  $\Pi_i(S) \subseteq M_i$ , so  $\Pi_1(S) \times \dots \times \Pi_n(S) \subseteq M_1 \times \dots \times M_n$ . Hence  $S \subseteq \Pi_1(S) \times \dots \times \Pi_n(S) \subseteq M_1 \times \dots \times M_n = \gamma(\langle M_1, \dots, M_n \rangle)$ . Furthermore  $\alpha(\gamma(\langle M_1, \dots, M_n \rangle)) \sqsubseteq \langle M_1, \dots, M_n \rangle$ , see exercise 8.3.2.

## Exercise 8.3

1. To prove monotonicity of  $\alpha$ , let  $S, S' \in D$  such that  $S \subseteq S'$ . Let  $x \in \Pi_i(S)$ , i. e., there exists  $\langle x_1, \dots, x_i, \dots, x_n \rangle \in S$  such that  $x = x_i$ . Then  $\langle x_1, \dots, x_i, \dots, x_n \rangle \in S'$ , so  $x = x_i \in \Pi_i(S')$ . Hence for all  $i$ ,  $\Pi_i(S) \subseteq \Pi_i(S')$ , and thus  $\alpha(S) = \langle \Pi_1(S), \dots, \Pi_n(S) \rangle \sqsubseteq \langle \Pi_1(S'), \dots, \Pi_n(S') \rangle = \alpha(S')$ .

To prove monotonicity of  $\gamma$ , let  $\langle M_1, \dots, M_n \rangle, \langle M'_1, \dots, M'_n \rangle \in D^\#$  such that  $\langle M_1, \dots, M_n \rangle \sqsubseteq \langle M'_1, \dots, M'_n \rangle$ . Then  $M_i \subseteq M'_i$  for all  $i$ , hence  $\gamma(\langle M_1, \dots, M_n \rangle) = M_1 \times \dots \times M_n \subseteq M'_1 \times \dots \times M'_n = \gamma(\langle M'_1, \dots, M'_n \rangle)$ .

2. To prove  $\text{id}_D \subseteq \gamma \circ \alpha$ , let  $S \in D$ . If  $\langle x_1, \dots, x_n \rangle \in S$  then  $x_i \in \Pi_i(S)$  for all  $i$ , so  $\langle x_1, \dots, x_n \rangle \in \Pi_1(S) \times \dots \times \Pi_n(S) = \gamma(\langle \Pi_1(S), \dots, \Pi_n(S) \rangle) = \gamma(\alpha(S))$ . Hence  $S \subseteq \gamma(\alpha(S))$ .

To prove  $\alpha \circ \gamma \sqsubseteq \text{id}_{D^\#}$ , let  $\langle M_1, \dots, M_n \rangle \in D^\#$ . Then  $\alpha(\gamma(\langle M_1, \dots, M_n \rangle)) = \alpha(M_1 \times \dots \times M_n) = \langle \Pi_1(M_1 \times \dots \times M_n), \dots, \Pi_n(M_1 \times \dots \times M_n) \rangle$ . Now we have to distinguish two cases. If  $M_1 \times \dots \times M_n = \emptyset$  then for all  $i$ ,  $\Pi_i(M_1 \times \dots \times M_n) = \emptyset$  otherwise  $\Pi_i(M_1 \times \dots \times M_n) = M_i$ . Thus,  $\alpha(\gamma(\langle M_1, \dots, M_n \rangle))$  either equals  $\langle M_1, \dots, M_n \rangle$  or  $\langle \emptyset, \dots, \emptyset \rangle$ . Hence in any case,  $\alpha(\gamma(\langle M_1, \dots, M_n \rangle)) \sqsubseteq \langle M_1, \dots, M_n \rangle$ .

3. Let  $S \in D$  and  $\langle M_1, \dots, M_n \rangle \in D^\#$ . Then

$$\begin{aligned}
S \subseteq \gamma(\langle M_1, \dots, M_n \rangle) &\Leftrightarrow S \subseteq M_1 \times \dots \times M_n \\
&\Leftrightarrow \forall i \in \{1, \dots, n\} \forall \langle x_1, \dots, x_i, \dots, x_n \rangle \in S : x_i \in M_i \\
&\Leftrightarrow \forall i \in \{1, \dots, n\} : \Pi_i(S) \subseteq M_i \\
&\Leftrightarrow \langle \Pi_1(S), \dots, \Pi_n(S) \rangle \sqsubseteq \langle M_1, \dots, M_n \rangle \\
&\Leftrightarrow \alpha(S) \sqsubseteq \langle M_1, \dots, M_n \rangle .
\end{aligned}$$

**Exercise 8.4** Set  $D^{\#'} = \{\langle \emptyset, \dots, \emptyset \rangle\} \cup \{\langle M_1, \dots, M_n \rangle \in 2^{D_1 \times \dots \times D_n} \mid M_i \neq \emptyset \text{ for all } i\}$ . Abstraction and concretization mappings are as before, i. e.,  $\alpha'(S) = \langle \Pi_1(S), \dots, \Pi_n(S) \rangle$  and  $\gamma'(\langle M_1, \dots, M_n \rangle) = M_1 \times \dots \times M_n$ , hence the properties demanded in exercises 8.2 and 8.3 continue to hold. To see why  $\alpha' \circ \gamma' = \text{id}_{D^{\#'}}$  recall the case distinction in the proof of  $\alpha \circ \gamma \sqsubseteq \text{id}_{D^\#}$ .