

## Polymorphic Lambda Calculus F

### Term-Type Presentation

$X \in TVar$

$T \in Ty = X \mid T \rightarrow T \mid \forall X.T$

$x \in Var$

$t \in Ter = x \mid \lambda x:T.t \mid tt \mid \lambda X.t \mid tT$

$\Gamma \in Var \rightarrow Ty$

$$\frac{\Gamma x = T}{\Gamma \vdash x : T} \quad \frac{\Gamma[x := T] \vdash t : T'}{\Gamma \vdash \lambda x:T.t : T \rightarrow T'} \quad \frac{\Gamma \vdash t_1 : T \rightarrow T' \quad \Gamma \vdash t_2 : T}{\Gamma \vdash t_1 t_2 : T'}$$

$$\frac{\Gamma \vdash t : T \quad X \text{ not free in } \Gamma}{\Gamma \vdash \lambda X.t : \forall X.T} \quad \frac{\Gamma \vdash t : \forall X.T'}{\Gamma \vdash tT : T'[X := T]}$$

### Uniform Presentation

$x \in Var$

$e \in Exp = x \mid \lambda x:e.e \mid ee \mid \Pi x:e.e \mid \star \mid \square$

$\Gamma \in Var \rightarrow Exp$

$\forall X.T \rightsquigarrow \Pi X:\star.T$

$T \rightarrow T' \rightsquigarrow \Pi X:T.T'$

$\lambda X.t \rightsquigarrow \lambda X:\star.t$

Define **sorting relation**  $\Gamma \vdash e : e'$  such that:

- **$e$  type**  $\iff \exists \Gamma: \Gamma \vdash e : \star$
- **$e$  term**  $\iff \exists \Gamma, e': \Gamma \vdash e : e' \wedge \Gamma \vdash e' : \star$
- **Sort Hierarchy** term : type :  $\star$  :  $\square$

Will use **suggestive Metavariables**:

- $e$  for terms and types
- $s$  for types and  $\star$  (sorts)
- $k$  for  $\star$  and  $\square$  (kinds)

$\Gamma, x : s \stackrel{\text{def}}{=} \Gamma \cup \{(x, s)\}$  where  $x \notin \text{Dom } \Gamma$

$$\begin{array}{c}
\frac{}{\Gamma, x : s \vdash x : s} \quad \frac{}{\Gamma \vdash \star : \square} \\
\frac{\Gamma \vdash s : k \quad k \in \{\star, \square\} \quad \Gamma, x : s \vdash s' : \star}{\Gamma \vdash \Pi x : s. s' : \star} \\
\frac{\Gamma \vdash e_1 : \Pi x : s. s' \quad \Gamma \vdash e_2 : s}{\Gamma \vdash e_1 e_2 : s' [x := e_2]} \\
\frac{\Gamma \vdash s : k \quad \Gamma, x : s \vdash e : s' \quad s'' = \Pi x : s. s' \quad \Gamma \vdash s'' : k'}{\Gamma \vdash \lambda x : s. e : s''}
\end{array}$$

$\Gamma$  well-formed:

- $\emptyset$  well-formed
- if  $\Gamma$  well-formed and  $\Gamma \vdash s : k$ , then  $\Gamma, x : s$  well-formed

### Uniform de Bruijn Presentation

$$n \in \text{Var} = \mathbb{N}$$

$$e \in \text{Exp} = n \mid \lambda e. e \mid e e \mid \Pi e. e \mid \star \mid \square$$

$$\Gamma \in \mathcal{L}(\text{Exp})$$

$$\text{Notation: } \Gamma = [e_1, \dots, e_n], \quad \Gamma.i = e_i$$

$$\begin{array}{c}
\frac{s = \text{up}(n+1)(\Gamma.(n+1))}{\Gamma \vdash n : s} \quad \frac{}{\Gamma \vdash \star : \square} \\
\frac{\Gamma \vdash s : k \quad k \in \{\star, \square\} \quad s :: \Gamma \vdash s' : \star}{\Gamma \vdash \Pi s. s' : \star} \\
\frac{\Gamma \vdash e_1 : \Pi s. s' \quad \Gamma \vdash e_2 : s}{\Gamma \vdash e_1 e_2 : \beta s' e_2} \\
\frac{\Gamma \vdash s : k \quad s :: \Gamma \vdash e : s' \quad s'' = \Pi s. s' \quad \Gamma \vdash s'' : k'}{\Gamma \vdash \lambda s. e : s''}
\end{array}$$