

Subtyping

A **subtype relation** is a partial order $\leq$ on the set of types that satisfies the **subsumption property**

$$\Gamma \vdash t : T \wedge \Gamma' \leq \Gamma \wedge T \leq T' \Rightarrow \Gamma' \vdash t : T'$$

where

$$\Gamma \leq \Gamma' \stackrel{\text{def}}{\iff} \text{Dom } \Gamma = \text{Dom } \Gamma' \wedge \forall X \in \text{Dom } \Gamma: \Gamma X \leq \Gamma' X$$

Type systems with a subtype relation generalize ordinary type systems, which are obtained by taking the identity relation as subtype relation. In place of the Unique Type Property we now require the more general **Least Type Property**:

$$\{ T \mid \Gamma \vdash t : T \} \neq \emptyset \Rightarrow \{ T \mid \Gamma \vdash t : T \} \text{ has a least element}$$

As always, we are only interested in typing relations that have the Preservation and the Progress Property. The **strong typing relation** $\vdash$ is characterized as follows:

$$\Gamma \vdash t : T \iff T \text{ is the least element of } \{ T' \mid \Gamma \vdash t : T' \}$$

Proposition 1 $\Gamma \vdash t : T \iff \exists T': \Gamma \vdash t : T' \wedge T' \leq T$

STR

STR is a simply typed lambda calculus with records and subtyping defined as follows:

$$l \in \text{Lab}$$

$$F \in \text{Lab} \xrightarrow{\text{fin}} \text{Ty}$$

$$T \in \text{Ty} = \text{Top} \mid T \rightarrow T \mid \Pi F$$

$$x \in \text{Var}$$

$$f \in \text{Lab} \xrightarrow{\text{fin}} \text{Ter}$$

$$t \in \text{Ter} = x \mid \lambda x:T.t \mid t t \mid \pi f \mid t.l$$

$$\Gamma \in \text{Var} \rightarrow \text{Ty}$$

$$(\lambda x:T.t)t' \rightarrow t[x := t']$$

$$\pi f.l \rightarrow fl \quad \text{if } l \in \text{Dom } f$$

$$\begin{array}{c}
\frac{}{T \leq \text{Top}} \quad \frac{T'_1 \leq T_1 \quad T_2 \leq T'_2}{T_1 \rightarrow T_2 \leq T'_1 \rightarrow T'_2} \quad \frac{\text{Dom } F' \subseteq \text{Dom } F \quad \forall l \in \text{Dom } F': Fl \leq F'l}{\Pi F \leq \Pi F'} \\
\\
\frac{\Gamma x = T}{\Gamma \vdash x : T} \quad \frac{\Gamma[x := T] \vdash t : T'}{\Gamma \vdash \lambda x:T. t : T \rightarrow T'} \quad \frac{\Gamma \vdash t : T'' \rightarrow T \quad \Gamma \vdash t' : T' \quad T' \leq T''}{\Gamma \vdash t t' : T} \\
\\
\frac{\text{Dom } f = \text{Dom } F \quad \forall l \in \text{Dom } f: \Gamma \vdash fl : Fl}{\Gamma \vdash \pi f : \Pi F} \quad \frac{\Gamma \vdash t : \Pi F \quad Fl = T}{\Gamma \vdash t.l : T} \\
\\
\Gamma \vdash t : T \stackrel{\text{def}}{\iff} \exists T': \Gamma \vdash t : T' \wedge T' \leq T
\end{array}$$

Note that we have defined $\vdash$ before $\vdash$. The following properties can be shown:

- $\leq$ is a partial order (by induction on types)
- $\vdash$ has the Unique Type Property (by induction on types).
- $\vdash$ has the Least Type Property (obvious)
- $\leq$ and $\vdash$ have the Subsumption Property (by induction on terms).
- Preservation, Progress and Termination.

Proposition 2 (Lubs and Glbs)

1. Two types always have a least upper bound.
2. If two types have a lower bound, they have a greatest lower bound.