



## Assignment 3 Semantics, WS 2009/10

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[www.ps.uni-sb.de/courses/sem-ws09/](http://www.ps.uni-sb.de/courses/sem-ws09/)

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Hand in by 11.59am, Tuesday, November 10

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There is a Standard ML file for this assignment. Send your solutions to Exercise 3.3 in a file named `lastname.sm1` to [doczka1@ps.uni-sb.de](mailto:doczka1@ps.uni-sb.de), and make sure that the entire file compiles without errors.

**Exercise 3.1 (System T Shallow)** Here is a shallow implementation of System T in Standard ML:

```
datatype nat = 0 | S of nat
fun natrec 0      x _ = x
  | natrec (S n) x f = f n (natrec n x f)
```

Write the following procedures using only *natrec* for recursion.

- a) *iszero* : *nat* → *bool*
- b) *pred* : *nat* → *nat*
- c) *add* : *nat* → *nat* → *nat*
- d) *mul* : *nat* → *nat* → *nat*
- e) *fac* : *nat* → *nat*

**Exercise 3.2 (PCF<sup>-</sup> Shallow)** Here is a shallow implementation of PCF<sup>-</sup> in Standard ML:

```
datatype nat = 0 | S of nat
fun natcase 0 x _ = x
  | natcase (S n) x f = f n
fun fix f x = f (fix f) x
```

A PCF<sup>-</sup> procedure for addition looks as follows:

```
val add = fix(fn f => fn x => fn y => natcase x y (fn x' => f x' (S y)))
```

Write and test PCF<sup>-</sup> procedures for multiplication and factorial.

**Exercise 3.3 (PCF<sup>-</sup> Deep)** We implement the abstract syntax of PCF<sup>-</sup> in Standard ML as follows:

```
datatype ty = Nat | P of ty * ty
type var = string
datatype ter = V of var | A of ter * ter | L of var * ty * ter
  | 0 | S of ter | C of ter * ter * ter | F of ter
```

- a) Write a procedure  $isVal : ter \rightarrow bool$  that tests whether a term is a value.
- b) Write a procedure  $elab : (var \rightarrow ty) \rightarrow ter \rightarrow ty$  that yields the type of a well typed term. Raise the exception *Error* if the term is not well-typed. Implement type environments as follows:  

```
exception Error
fun empty x = raise Error
fun update f x a y = if y=x then a else f y
```
- c) Write a procedure  $subst : var \rightarrow ter \rightarrow ter \rightarrow ter$  that computes  $[x := s]t$  if  $s$  is closed.
- d) Write a procedure  $eval : ter \rightarrow ter$  that yields the value of a closed term if it exists. Raise the exception *Error* if  $eval$  must quit because of a type inconsistency or a free variable occurrence.