

Assignment 13 Semantics, WS 2011-2012

Prof. Dr. Gert Smolka, Dr. Chad Brown www.ps.uni-saarland.de/courses/sem-ws11/

Read Chapters 7 and 8 of the Lecture Notes

Note: This assignment is relevant for the Endterm.

Exercise 13.1 Give an example of a closed term t in the simply-typed lambda calculus such that there is no type T such that $R_T t$. Give an example of such a t that also terminates (relative to the nondeterministic weak reduction in Chapter 7 of the lecture notes).

Exercise 13.2 Consider the simply typed lambda calculus with the typing relation, the nondeterministic weak reduction, and the logical relation *R* in Chapter 7 of the lecture notes. Which of the following statements are true?

- a) If $R_T t$, then $\emptyset \vdash t : T$.
- b) If $R_T t$, then t terminates.
- c) If $\emptyset \vdash t : T$, then $R_T t$.
- d) If $\emptyset \vdash t : T$, then t terminates.
- e) If $\Gamma \vdash t : T$ and θ is a closed substitution with the same domain as Γ , then $\emptyset \vdash \theta t : T$.
- f) If $\Gamma \vdash t : T$ and $R_{\Gamma}\theta$, then $R_{T}(\theta t)$.
- g) If $t \Rightarrow t'$ and $R_T t$, then $R_T t'$.
- h) If $t \Rightarrow t'$ and $R_T t'$, then $R_T t$.

Remark: Exercises 13.3 - 13.9 concern the Calculus of Constructions (Chapter 8 of the Lecture Notes).

Exercise 13.3 Suppose $\emptyset \vdash s : t$ and s is normal. Find out whether s can be a variable or an application.

Exercise 13.4 Suppose the typing $\Gamma \vdash s : t$ is derivable and t reduces to u. Explain why the typing $\Gamma \vdash s : u$ is derivable.

Exercise 13.5 You can experiment with the typing rules in Coq. Do the following examples by hand (taking Type to be U_0) and check your results with Coq.

```
Check fun (s : Type) (t : s \rightarrow Type) => forall x : s, t x.

Check fun (s u : Type) (t : s \rightarrow u) => fun x : s => t x.

Check fun (u : Type) (v : u \rightarrow Type) (s : forall x : u, v x) (t : u) => s t.

Check fun X : Type => X \rightarrow forall X : Type, X.
```

Exercise 13.6 Derive the typing

$$\emptyset \vdash \lambda x : U_0.\lambda x : x.x : \forall x : U_0. \forall y : x.x$$

Exercise 13.7 Derive the following typings.

- a) $X: U_0 \vdash (\lambda Y: U_0.Y)X: U_0$
- b) $X: U_0 \vdash (\lambda X: U_0.X)X: U_0$

Exercise 13.8 Determine normal types of the following terms and check your results with Coq.

- a) $\forall x : U_0.x$
- b) $\lambda x: U_0. \forall y: U_0.x \rightarrow y$
- c) $\lambda f: U_0 \rightarrow U_1. \forall x: U_0. f x$
- d) $\lambda xyz:U_0.\lambda f:x \to y.\lambda g:y \to z.\lambda w:x.g(fw)$

Exercise 13.9 Convince yourself that the following terms are ill-typed.

- a) $\forall x : U_0 . \forall y : x . y$
- b) $\lambda f: U_1 \rightarrow U_0. \forall x: f U_0.x$
- c) $\lambda x y z : U_0.\lambda f : x \to y.\lambda g : y \to z.$ $\forall p : (x \to z) \to U_0. \ p(\lambda w : x.g(f w)) \to p(\lambda w : x.w)$