



## Assignment 2 Semantics, WS 2013/14

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[www.ps.uni-saarland.de/courses/sem-ws13/](http://www.ps.uni-saarland.de/courses/sem-ws13/)

**Definitions:**  $I := \lambda x.x$ ,  $B := \lambda f g x.f(gx)$ ,  $\hat{n} := \lambda f x.f^n x$ ,  $\text{succ} := \lambda n f x.f(nfx)$ .

**Exercise 2.1** Prove the following equivalences by hand:

- |   |  |
|---|--|
| a) $\hat{n} \equiv \lambda f x. \hat{n} f x$  | f) $\widehat{S}n s \equiv B s (\hat{n} s)$                           |
| b) $\hat{n} \equiv \lambda f. \hat{n} f$      | g) $\hat{n} s \equiv \hat{n} (B s) I$                                |
| c) $\text{succ } s t \equiv B t (s t)$        | h) $\widehat{m + n} \equiv \lambda s. B (\widehat{m} s) (\hat{n} s)$ |
| d) $\widehat{S}n \equiv \text{succ } \hat{n}$ | i) $\widehat{m \cdot n} \equiv B \widehat{m} \hat{n}$                |
| e) $\widehat{S}n s t \equiv s (\hat{n} s t)$  | j) $\widehat{m^n} \equiv \hat{n} (B \widehat{m}) \hat{1}$            |

**Exercise 2.2** Define the following functions on Church numerals. You may use primitive recursion (prec).

a) Tetration:  ${}^n a = \underbrace{a^{a^{\dots^a}}}_n$ . We leave the value of  ${}^0 a$  unspecified.

b) Ackermann's function:

$$A(x, y) = \begin{cases} y + 1 & \text{if } x = 0 \\ A(x - 1, 1) & \text{if } y = 0 \\ A(x - 1, A(x, y - 1)) & \text{otherwise} \end{cases}$$

If you run into difficulties, read Chapter 4 in the ICL lecture notes.

**Exercise 2.3** Consider the following two definitions of multiplication.

$$\text{add} := \lambda m n f x. m f(nfx)$$

$$\text{mul}_1 := \lambda m n f. m(nf)$$

$$\text{mul}_2 := \lambda m n. m(\text{add } n)(\lambda f x. x)$$

The definitions agree on Church numerals, but are not equivalent. Find terms  $a_1, \dots, a_k$  such that:

$$\text{mul}_1 a_1 \dots a_k \succ^* \lambda x y. x$$

$$\text{mul}_2 a_1 \dots a_k \succ^* \lambda x y. y$$

**Exercise 2.4** Let  $s$  be an abstraction such that  $x$  is not free in  $s$ . Argue that  $\lambda x. s x \equiv s$ .

**Exercise 2.5 (Coq)** Show that the Church-Rosser Property (CRP) implies uniqueness of normal forms.