Universität des Saarlandes Programming Systems Lab Polymorphic Lambda Calculus with Dynamic Types

> Bachelor Thesis Task Formulation

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Motivation

Open Programming System:

- Not all components available at compile time
- Components linked dynamically
- Dynamic type checking needed

Possible construct: type case

- Notation: case $t_1:T_1$ of $x:T_2 \Rightarrow t_2$ else t_3
- Provides branching dependent on type T_1 of subterm t_1
- Evaluates to $t_2[x := t_1]$ iff $T_1 = T_2$ dynamically otherwise to t_3

Type Case

Example:

```
rep = \lambda X.\lambda x : X.case \ x : X \text{ of } x' : bool \Rightarrow "bool"
else case x : X of x' : int \Rightarrow "int"
else "unknown"
```

Given a type and a term of this type *rep* returns a string representation of the type.

Problem: Type case destroys parametricity of type abstraction.

abstype Number = int
implementation: [...]

rep Number $n \longrightarrow^*$ "int"

Dynamic Type Name Generation

Solution: Generate new type names dynamically.

- Notation: new X = T in t
- Type name X can be used in t in place of T.
- Use global state for dynamically generated type names instead of coercions (Rossberg's approach).

Example:

new X = int in abstype Number = Ximplementation: [...]

 $rep Number n \longrightarrow^* "unknown"$

Laziness

Call-by-need extension for simply typed λ -calculus:

- Notation: lazy x = t in t'
- Variable x can be used in t' in place of t (similar to let)
- Evaluate t as late as possible.
- Use global state for modelling relationship between x and t.

Similar construct can be used to express lazy linking:

lazy < X, x > = < T, t > in t'

Task Formulation

Develop a calculus, λ_F^N , an extension of system F with constructs for type case and a binder for new type names.

- Use global state instead of coercions.
- Use evaluation contexts.
- Prove the unique type, the progress and the preservation property.

Extend the simply typed λ -calculus with call-by-need evaluation and give proofs for the safety properties.

Use this construct to model lazy linking in system F or in λ_F^N

Give a model implementation.

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