Syntactic Theory of Finitary Sets

Denis Müller

Saarland University

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Introduction

Hereditarily finitary sets consist of finitely many hereditarily finitary sets. Well-founded sets do not have cycles and can hence be represented by binary trees. Semantically, this is how we can map finitary sets in our

representation to ordinary sets:

$$\blacktriangleright \llbracket s.t \rrbracket = \{\llbracket s \rrbracket\} \cup \llbracket t \rrbracket$$

Functions on Trees

List of all elements of a tree: L $\emptyset = []$ L (s . t) = s :: L t

Append: (a) \emptyset (a) t = t (s.s') (a) t = s . s' (a) t

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Properties of Equivalence Relations

Important properties that should be satisfied:

• Del : s.s.t
$$\equiv$$
 s.t

Swap : s.t.u \equiv t.s.u

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We inductively define a first equivalence relation \equiv satisfying the previous conditions

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$$\overline{\emptyset \equiv \emptyset}$$

$$\overline{s.s.t \equiv s.t}$$

$$\overline{s.t.u \equiv t.s.u}$$

$$\underline{s \equiv t \quad s \neq \emptyset \neq t}$$

$$\overline{t \equiv s}$$

$$\underline{s \equiv t \quad t \equiv u}$$

$$\underline{s \equiv s' \quad t \equiv t'}$$

$$\overline{s.t \equiv s'.t'}$$

 \equiv is (basically) the least congruence satisfying Del and Swap.

Properties of \equiv

- Symmetry, Transitivity, as well as the Deletion, Swap and Composition Rule are trivially fulfilled by ≡.
- ► Reflexivity is also admissible: s = s is easily established by induction on s.

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Motivation

Wanted: normal form and normalizer η s.t.

- normal (η s)
- normal $s \rightarrow$ normal $t \rightarrow s \equiv t \rightarrow s = t$

The goal is to proof decidability of \equiv \equiv will be used to define set relations on trees.

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Lexicographic comparison

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Order ::= LE | EQ | GT

cmp : Tree \rightarrow Tree \rightarrow Order

cmp \emptyset \ \emptyset = EQ

cmp \emptyset \ t = LT

cmp s \emptyset = GT

cmp (s1.s2) (t1.t2) = match (cmp s1 t1) with

| EQ \rightarrow cmp (s2 t2)

| x \rightarrow x end
```

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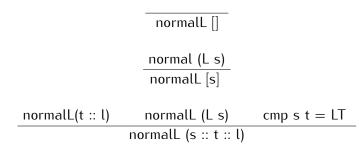
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Normal



normal is defined in terms of normalL as follows: normal s := normalL (L s).

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Insertion Sort

insert s
$$\emptyset$$
 = s. \emptyset
insert s (t1.t2) = match cmp s t1 with
|EQ \rightarrow t1.t2
|LT \rightarrow s.t1.t2
|GT \rightarrow t1 . insert s t2 end

sort $\emptyset = \emptyset$ sort (s.t) = insert (sort s) (sort t)

Properties of sort

We can prove the following properties by induction:

- s.t ≡ insert s t
- s ≡ sort s
- sort (s.s.t) = sort (s.t)
- sort(s.t.u) = sort(t.s.u)
- $s \equiv t \rightarrow sort s = sort t$
- normal (sort s)
- normal $s \rightarrow sort s = s (\rightarrow Idempotency of sort)$

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- two sorted sets are equivalent they are syntactically equal
- sort s ≡ s
- ► ≡ is transitive
- s = t is decidable, since finitary sets are a simple inductive type

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Element relationship

► $s \in t := t \equiv s.t$

$$\bullet \ s \in t \iff \exists t'. t' \in L t \land s \equiv t'.$$

Subset relationship

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Coinductive Bisimulation

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\begin{split} s &\sim t := \\ \forall \ s. \ s' &\in L(s) \rightarrow \exists \ t'. \ t \in L(t) \land s' \sim t' \land \\ \forall \ t. \ t' &\in L(t) \rightarrow \exists \ s'. \ s \in L(s) \land t' \sim s'. \end{split}
```

Reflexivity, Symmetry, Swap and Del can be proven easily.

Transitivity of \sim

 $\begin{array}{l} \mbox{Coinduction lemma needed:} \\ \mbox{Lemma TreeCoInd (R : Tree \rightarrow Tree \rightarrow Prop) :} \\ (\forall \ s \ t, \ R \ s \ t \rightarrow \\ (\forall \ s. \ s' \in L(s) \rightarrow $\exists \ t'. \ t \in L(t) \land $R \ s' \ t' \land \\ \forall \ t. \ t' \in L(t) \rightarrow $\exists \ s'. \ s \in L(s) \land $R \ t' \ s'.)) \rightarrow \\ \forall \ s \ t, \ R \ s \ t \rightarrow $s \sim t. \\ \end{array}$

Transitivity follows by using the lemma with R := (~ \circ ~) \Rightarrow s = t \rightarrow s ~ t If suffices to show: $s \sim t \rightarrow s \subseteq t \land t \subseteq s$ Stronger induction lemma:

 $\begin{array}{l} \text{Lemma TreeInduction } (\mathsf{P}:\mathsf{Tree}\to\mathsf{Prop}):\\ (\forall \ \mathsf{s} \ , (\forall \ \mathsf{t}, \ \mathsf{t}\in(\mathsf{L}\ \mathsf{s})\to\mathsf{P}\ \mathsf{t})\to\mathsf{P}\ \mathsf{s})\to\forall \mathsf{s}, \ \mathsf{P}\ \mathsf{s}. \end{array}$

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🛸 Sangiorgi, Davide. Introduction to bisimulation and coinduction. Cambridge University Press, 2011.