## Formalizing Stream-Calculus in Coq

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 $a_0 = 1$  $a_1 = 1$  $a_{n+2} = a_n + a_{n+1}$ 

## Fibonacci sequence as example

$$F_0 = 1$$
  

$$F_1 = 1$$
  

$$F_{n+2} = F_n + F_{n+1}$$

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$$F_0 = 1$$
  

$$F_1 = 1$$
  

$$F_{n+2} = F_n + F_{n+1}$$

$$\sum_{i=0}^{n} F_i = F_{n+2} - 1$$
$$F_n = \frac{1}{\sqrt{5}} (\phi^n - \hat{\phi}^n)$$

 $t_0, t_1, t_2, t_3, t_4...$ 

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#### Interpretation as Sequence

 $a_n := t_n$ 

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 $t_0, t_1, t_2, t_3, t_4...$ 

#### Interpretation as Sequence

 $a_n := t_n$ Interpretation as Stream head:  $s_0 := t_0$ tail:  $s' := t_1, t_2, t_3, t_4, \dots$ 

#### $1, 1, 1, 1, 1, \dots$

#### Interpretation as Sequence

 $a_n := 1$ 

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#### Interpretation as Sequence

 $a_0 := 1$  $a_{n+1} := a_n$ 

 $1, 1, 1, 1, 1, \dots$ 

#### Interpretation as Sequence

 $a_0 := 1$  $a_{n+1} := a_n$ 

## Interpretation as Stream head: $s_0 := 1$

tail: s' := s

 $1, 1, 1, 1, 1, \dots$ 

Interpretation as Sequence  $a_0 := 1$   $a_{n+1} := a_n$ Recursion Interpretation as Stream head:  $s_0 := 1$ tail: s' := sDifferential Equation

#### Adding Sequences $(a+b)_n := a_n + b_n$

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## Adding Streams $(r+s)_0 := r_0 + s_0$ (r+s)' := r' + s'

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Ascending Sequence  $a_0 := 0$  $a_{n+1} := a_n + 1$ 

Adding Sequences  $(a+b)_n := a_n + b_n$ 

Ascending Sequence  $a_0 := 0$  $a_{n+1} := a_n + 1$  Adding Streams  $(r + s)_0 := r_0 + s_0$ (r + s)' := r' + s'

Ascending Stream  $s_0 := 0$ s' := s + (1, 1, ...)

#### Some Definitions

$$\begin{array}{l} 0_0 := 0 \\ 0' := 0 \\ [t] := (t, 0, 0, \dots) \\ [t]_0 := t \\ [t]' := 0 \end{array}$$

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#### **Multiplying Streams**

$$(r \times s)_0 := r_0 \times s_0$$
  
 $(r \times s)' := [r_0] \times s' + r' \times s$ 

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#### Some Definitions

$$\begin{array}{l} 0_0 := 0 \\ 0' := 0 \\ [t] := (t, 0, 0, \dots) \\ [t]_0 := t \\ [t]' := 0 \end{array}$$

#### Behaviour of Multiplication

 $0 \times s = 0$ 

$$([t] imes s)_0 := t imes s_0$$
  
 $([t] imes s)' := [t] imes s' + 0 imes s$ 

 $[1] \times s = s$ 

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#### Behaviour of Multiplication

 $0 \times s = 0$ 

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 $([t] \times s)' := [t] \times s' + 0 \times s$ 

 $[1] \times s = s$ 

#### Multiplying Polynoms

$$a = a_0 x^0 + a_1 x^1 + \dots$$
  

$$b = b_0 x^0 + b_1 x^1 + \dots$$
  

$$a \times b = c_0 x^0 + c_1 x^1 + \dots$$

$$c_i = \sum_{l+k=i} a_l \times b_k$$
  
=  $a_0 \times b_i + \sum_{l+k=i-1} a_{l+1} \times b_k$ 

## Backshift X

$$egin{aligned} X &= (0,1,0,0,\dots) \ X_0 &:= 0 \ X' &:= [1] \end{aligned}$$

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$$(X \times s)_0 = 0$$
  
 $(X \times s)' = X_0 \times s' + X' \times s = s$ 

$$s = [s_0] + X \times s'$$

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# Dividing Streams $(s \times s^{-1}) = 1$

#### **Dividing Streams**

$$(s \times s^{-1}) = 1$$

► 
$$1 = (s \times s^{-1})_0 = s_0 \times s_0^{-1}$$
  
 $\Rightarrow s_0^{-1} = 1/s_0$ 

#### **Dividing Streams**

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**Dividing Sequences** 

$$s_n^{-1} = \frac{-\sum_{l+k=n-1} s_{l+1} \times s_k^{-1}}{s_0}$$

#### **Dividing Streams**

$$(s \times s^{-1}) = 1$$

**Dividing Sequences** 

$$s_n^{-1} = \frac{-\sum_{l+k=n-1} s_{l+1} \times s_k^{-1}}{s_0}$$

Constant Stream by Division

$$(1-X)^{-1} = (1,1,1,\dots)$$

## Ring on Streams

#### Addition

$$(a+b)_0 := a_0 + b_0$$
  
 $(a+b)' := a' + b'$ 

#### Subtraction

$$(-a)_0 := -a_0$$
  
 $(-a)' := -a'$ 

#### Multiplication

#### Division

$$(s^{-1})_0 := s_0^{-1} \ (s^{-1})' := -s' \times ([s_0] \times s)^{-1}$$

Zero- and One-Elements  $\begin{array}{l} 0_0:=0\\ 0':=0\\ 1:=[1] \end{array}$ 

## Equalities on Streams

Pointwise Equality

 $a = b :\leftrightarrow a_0 = b_0 \wedge a' = b'$ 

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Equality up-to

$$a =_{n+1} b : \leftrightarrow a_0 = b_0 \land a' =_n b'$$
  
 $\forall a, b : a =_0 b$ 

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Pointwise Equality

$$a = b : \leftrightarrow a_0 = b_0 \land a' = b'$$

Equality up-to

$$a =_{n+1} b : \leftrightarrow a_0 = b_0 \wedge a' =_n b' \ orall a, b : a =_0 b$$

$$(\forall n: a =_n b) \leftrightarrow a = b$$

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Fibonacci as Sequence

 $F_0 := 1$  $F_1 := 1$  $F_{n+2} := F_n + F_{n+1}$ 

Fibonacci as Sequence

 $F_0 := 1$  $F_1 := 1$  $F_{n+2} := F_n + F_{n+1}$  Fibonacci as Stream

 $\begin{array}{l} \textit{fib}_0 := 1 \\ \textit{fib}_1 := 1 \\ \textit{fib}'' := \textit{fib} + \textit{fib}' \end{array}$ 

Fibonacci as Sequence

 $F_0 := 1$  $F_1 := 1$  $F_{n+2} := F_n + F_{n+1}$ 

#### Fibonacci as Stream

 $\begin{array}{l} \textit{fib}_0 := 1 \\ \textit{fib}_1 := 1 \\ \textit{fib}'' := \textit{fib} + \textit{fib}' \end{array}$ 

$$\begin{aligned} fib &= fib_0 + X \times fib' \\ &= fib_0 + X \times (fib_1 + X \times fib'') \\ &= 1 + X \times 1 + X^2 \times fib'' \\ &= 1 + X \times 1 + X^2 \times (fib + fib') \\ &= 1 + X \times (1 + X \times fib') + X^2 \times fib \\ &= 1 + X \times fib + X^2 \times fib \end{aligned}$$

Fibonacci as Sequence

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Fibonacci closed

$$fib = \frac{1}{1-X-X^2}$$

## Example of reasoning with Streams

#### Fibonacci closed

$$fib = \frac{1}{1-X-X^2}$$

$$fib' = \frac{1+X}{1-X-X^2}$$

#### Fibonacci closed

$$fib = \frac{1}{1-X-X^2}$$

$$\textit{fib}' = \frac{1+X}{1-X-X^2}$$

$$\sum_{k=0}^{n} F_k$$

$$rac{1}{1-X} imesrac{1}{1-X-X^2}$$

#### Fibonacci closed

$$fib = \frac{1}{1-X-X^2}$$

$$fib' = \frac{1+X}{1-X-X^2}$$

$$\sum_{k=0}^{n} F_k$$

$$\frac{1}{1-X} imes \frac{1}{1-X-X^2} = \frac{A}{1-X} + \frac{B}{1-X-X^2}$$

#### Fibonacci closed

$$fib = \frac{1}{1-X-X^2}$$

$$fib' = \frac{1+X}{1-X-X^2}$$

$$\sum_{k=0}^{n} F_k$$

$$\frac{1}{1-X} \times \frac{1}{1-X-X^2} = \frac{-1}{1-X} + \frac{2+X}{1-X-X^2}$$

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$$fib = \frac{1}{1-X-X^2}$$

$$fib' = \frac{1+X}{1-X-X^2}$$

$$\sum_{k=0}^{n} F_k$$

$$\frac{1}{1-X} \times \frac{1}{1-X-X^2} = \frac{-1}{1-X} + \frac{2+X}{1-X-X^2}$$
$$= -\frac{1}{1-X} + fib''$$

#### Fibonacci closed

$$fib = \frac{1}{1-X-X^2}$$

$$\textit{fib}' = \frac{1+X}{1-X-X^2}$$

$$\sum_{k=0}^{n} F_k = F_{n+2} - 1$$

$$\frac{1}{1-X} \times \frac{1}{1-X-X^2} = \frac{-1}{1-X} + \frac{2+X}{1-X-X^2}$$

$$=-rac{1}{1-X}+fib''$$

### Implementing Streams in Coq

Stream  $T := \mathbb{N} \to T$ head s := (s 0)tail  $s := \lambda n, s(n + 1)$ 

## Implementing Multiplication

#### Definition

Implementation

$$a \times b := \lambda n, \text{match } n \text{ with}$$
$$|0 \Rightarrow a \, 0 \times b \, 0$$
$$|Sn' \Rightarrow (a_0 \times b' + a' \times b) \, n'$$

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## Implementing Multiplication

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## Implementing Division

### Definition

$$(s^{-1})_0 := s_0^{-1} \ (s^{-1})' := -s' imes ([s_0] imes s)^{-1}$$

Implementation

$$inv a := \lambda n$$
,match  $n$  with  
 $|0 \Rightarrow inv(a 0)$   
 $|Sn' \Rightarrow (-a' \times inv([a_0] \times a)) n'$ 

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tail: 
$$\underbrace{Stream}_{input} \rightarrow Stream$$

$$inv \ a := \lambda n, \text{match } n \text{ with}$$

$$|0 \Rightarrow inv(a \ 0)$$

$$|Sn' \Rightarrow (-a' \times inv([a_0] \times a)) \ n'$$
pre-tail: 
$$\underbrace{Stream}_{\text{input}} \rightarrow \underbrace{(Stream \rightarrow Stream)}_{\text{inverse}} \rightarrow Stream$$

tail: 
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pre-tail: 
$$\underbrace{Stream}_{input} \rightarrow \underbrace{(Stream \rightarrow Stream)}_{inverse} \rightarrow Stream$$
  
(pre\_inv\_tail s :=  $\lambda inv, -s' \times inv([s_0] \times s)$ 

$$\textbf{tail:} \; \underbrace{\textit{Stream}}_{\text{input}} \rightarrow \textit{Stream}$$

pre-tail: 
$$\underbrace{Stream}_{input} \rightarrow \underbrace{(Stream \rightarrow Stream)}_{approximization} \rightarrow Stream$$
  
(pre\_inv\_tail s :=  $\lambda inv, -s' \times inv([s_0] \times s)$ 

tail: 
$$\underbrace{Stream}_{input} \rightarrow Stream$$

$$inv a := \lambda n$$
,match  $n$  with  
 $|0 \Rightarrow inv(a 0)$   
 $|Sn' \Rightarrow (-a' \times inv([a_0] \times a)) n'$ 

pre-tail:  $\underbrace{Stream}_{input} \rightarrow \underbrace{(Stream \rightarrow Stream)}_{approximization} \rightarrow Stream$ (pre\_inv\_tail s :=  $\lambda inv, -s' \times inv([s_0] \times s)$ Causality

 $causal \phi :\Leftrightarrow \forall n, a1, a2 : (\forall s : a1 s =_n a2 s) \rightarrow \phi a1 =_n \phi a2$ 

tail: 
$$\underbrace{Stream}_{input} \rightarrow Stream$$

$$inv a := \lambda n$$
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pre-tail:  $\underbrace{Stream}_{input} \rightarrow \underbrace{(Stream \rightarrow Stream)}_{approximization} \rightarrow Stream$ (pre\_inv\_tail  $s := \lambda inv, -s' \times inv([s_0] \times s)$ Causality causal  $\phi :\Leftrightarrow \forall n, a1, a2 : (\forall s : a1 s =_n a2 s) \rightarrow \phi a1 =_n \phi a2$ 

 $C Stream := \{\phi : (Stream \rightarrow Stream) \rightarrow Stream | causal \phi\}$ 

tail: 
$$\underbrace{Stream}_{input} \rightarrow Stream$$

$$inv a := \lambda n, \text{match } n \text{ with}$$
$$|0 \Rightarrow inv(a 0)$$
$$|Sn' \Rightarrow (-a' \times inv([a_0] \times a)) n'$$

pre-tail: 
$$\underbrace{Stream}_{input} \rightarrow C Stream$$
  
(pre\_inv\_tail  $s := \lambda inv, -s' \times inv([s_0] \times s)$   
Causality

 $causal \phi :\Leftrightarrow \forall n, a1, a2 : (\forall s : a1 s =_n a2 s) \to \phi a1 =_n \phi a2$ 

C Stream := { $\phi$  : (Stream  $\rightarrow$  Stream)  $\rightarrow$  Stream|causal  $\phi$ } By this: stepwise approximation of  $s^{-1}$ . Implementing Division (cnt.)

### Definition

$$egin{array}{l} (s^{-1})_0 := s_0^{-1} \ (s^{-1})' := -s' imes ([s_0] imes s)^{-1} \end{array}$$

Implementation

 $inv\_head a := inv(a0)$   $pre\_inv\_tail := \lambda inv\_app, -a' \times inv\_app([a_0] \times a)$  $inv := approx\_stream inv\_head pre\_inv\_tail$ 

### Approximazation revised

### Causality

 $causal \phi :\Leftrightarrow \forall n, a1, a2 : (\forall s : a1 s =_n a2 s) \rightarrow \phi a1 =_n \phi a2$ 

$$C Stream := \{\phi : (Stream \rightarrow Stream) \rightarrow Stream | causal \phi\}$$
  
pre-tail: 
$$\underbrace{Stream}_{input} \rightarrow C Stream$$

### Approximazation revised

### Causality

 $causal \phi :\Leftrightarrow \forall n, a1, a2 : (\forall s : a1 s =_n a2 s) \rightarrow \phi a1 =_n \phi a2$ 

$$CX := \{\phi : (X \to Stream) \to Stream | causal \phi\}$$
  
pre-tail:  $X \to CX$   
input

### Approximazation revised

### Causality

 $\textit{causal } \phi :\Leftrightarrow \forall n, a1, a2 : (\forall s : a1 \, s =_n a2 \, s) \rightarrow \phi \, a1 =_n \phi \, a2$ 

$$CX := \{\phi : (X \to Stream) \to Stream | causal \phi\}$$
  
pre-tail:  $X \to CX$ 

Way to solve a system of causal differential equations by approximazation.

## Summary

- Streams as simple and clean representation of sequences
- A way to workaround Coq's restrictions with respect to coinductive definitions.
- A way to solve a system of causal differential equations

## Further Reading

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- Herbert S. Wilf. generatingfunctionology. Academic Press, 1990.

Definition Pointwise Equality  $a = b : \leftrightarrow a_0 = b_0 \land a' = b'$ 

### Definition Pointwise Equality $a = b : \leftrightarrow a_0 = b_0 \land a' = b'$

$$a \equiv b : \leftrightarrow a_0 \equiv b_0 \land a \equiv$$

### Implementation

$$a = b := \forall n, a n = b n$$

Definition Equality up-to

$$a =_{n+1} b : \leftrightarrow a_0 = b_0 \land a' =_n b'$$
  
 $\forall a, b : a =_0 b$ 

### **Definition Pointwise Equality**

$$a=b: \leftrightarrow a_0=b_0 \wedge a'=b'$$

Implementation

$$a = b := \forall n, a n = b n$$

Definition Pointwise Equality  

$$a = b :\leftrightarrow a_0 = b_0 \land a' = b'$$
  
Implementation  
 $a = b := \forall n, a n = b n$ 

# Definition Equality up-to $a =_{n+1} b : \leftrightarrow a_0 = b_0 \land a' =_n b'$ $\forall a, b : a =_0 b$

Implementation

$$r =_0 s$$

$$\frac{r_0 = s_0 \qquad r' =_n s'}{r =_{n+1} s}$$

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## Implementing Addition

### Implementation

### Definition

$$(a+b)_0 := a_0 + b_0$$
  
 $(a+b)' := a' + b'$ 

$$a + b := \lambda n$$
,match  $n$  with  
 $|0 \Rightarrow a 0 + b 0$   
 $|Sn' \Rightarrow (a' + b')n'$ 

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### Definition

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### Implementation

$$a+b := \lambda n, a n + b n$$