

Verification of a Case Analysis Plugin in MetaCoq

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13. März 2020



METACOQ

- framework implementing Coq's logic
- monadic manipulation for plugins
- TemplateCoq implementation: bare type system
- pCUIC implementation: more abstract with clear semantic
- typing is provable
- verification using meta proofs

METACOQ

```

[]
ind_finite   := Finite;
ind_params  := 0;
ind_univers := [];
ind_universes := Nonomorphic_cts (LevelSetProp.of_list [], ConstraintSet.empty) []
[]
wind_entry_record := None;
wind_entry_finite := Finite;
wind_entry_params := [];
wind_entry_universes := Nonomorphic_cts
  (LevelSet.this := [], LevelSet.is_ok := LevelSet.Raw.empty_ok [],
   [ConstraintSet.this := [], ConstraintSet.is_ok := ConstraintSet.Raw.empty_ok []]);
wind_entry_private := None
  
```

Transform

```

(tlambd (nName "p")
  (tProd (nName "inst")
    (tInd [] inductive_mind := "Coq.Init.Datatypes.nat"; inductive_ind := 0 [] [])
    (tSort (NEL.sing (Level.Level "Top.140", false))))))
(tlambd (nName "H_0")
  (tConstruct [] inductive_mind := "Coq.Init.Datatypes.nat"; inductive_ind := 0 [] 0 [])
  (tlambd (nName "H_S")
    (tProd nAnon (tInd [] inductive_mind := "Coq.Init.Datatypes.nat"; inductive_ind := 0 [] 1)
      (tApp (tRel 2)
        [tApp
          [tConstruct [] inductive_mind := "Coq.Init.Datatypes.nat"; inductive_ind := 0 [] 1
            [] (tRel 0)]]])
        [dName := nName "??";
         dType := tProd (nName "inst")
           [tInd [] inductive_mind := "Coq.Init.Datatypes.nat"; inductive_ind := 0 [] 0]
           (tApp (tRel 3) (tRel 0));
          dBody := tlambd (nName "inst")
            [tInd [] inductive_mind := "Coq.Init.Datatypes.nat"; inductive_ind := 0 [] 0]
            (tCase
              [{} inductive_mind := "Coq.Init.Datatypes.nat"; inductive_ind := 0 {} 0,
               (tlambd (nName "inst")
                 (tInd
                   [{} inductive_mind := "Coq.Init.Datatypes.nat"; inductive_ind := 0 []
                     [] (tApp (tRel 5) (tRel 0))] (tRel 0)
                   [{} (0, tRel 3);
                    {} (1,
                      tlambd nAnon
                        (tInd
                          [{} inductive_mind := "Coq.Init.Datatypes.nat"; inductive_ind := 0 []
                            [] (tApp (tRel 3) (tRel 0))]))];
                    fArg := 0 [] 0)]]])
                ]))
          ]))
  )
  
```

Quote

Inductive N : Set := 0 : N | S : N → N

Unquote

```

λ (p : N → Type) (H_0 : p 0) (H_S : ∀ H : N, p (S H)) ⇒
fix f (inst : N) : p inst :=
  match inst as inst₀ return (p inst₀) with
  | 0 ⇒ H_0
  | S x ⇒ H_S x
  end
: ∀ p : N → Type, p 0 → (∀ H : N, p (S H)) → ∀ inst : N, p inst
  
```

METACOQ

typing $\Sigma; \Gamma \vdash t : T$

Inductive typing $\Sigma \Gamma : \text{term} \rightarrow \text{term} \rightarrow \mathbb{T} :=$

| type_Rel : ...

| type_Sort : ...

| type_Prod : ...

| type_Lambda : ...

| type_App : ...

| type_Ind : ...

| type_Construct : ...

| type_Case : ...

| type_Fix : ...

| type_Conv : ...

| ...

typingSpine $\Sigma; \Gamma \vdash t^T xs : T_2$



CHANGES

- changed plugin to use predefined functions
- implementation of induction
- plugin simplification

CORRECTNESS

```
wf  $\Sigma \Gamma \rightarrow$   
on_ind_body mind mind_body ind ind_body  $\rightarrow$   
declared_inductive  $\Sigma$  mind_body ind ind_body  $\rightarrow$   
 $\forall T t \text{ name},$   
createElim ind = Some (t,name)  $\rightarrow$   
createElimType ind = Some T  $\rightarrow$   
 $\Sigma; \Gamma \vdash t : T.$ 
```

- concrete type to avoid existentials



PROOF STRUCTURE

1. intros
2. typing of p
3. intros all cases
4. resolve fixpoint
5. wf_local of environment
6. transform to lambda and type to prod
7. intros
8. typing of instance
9. simplification
10. η -Conversion
11. case analysis typing

```

E≤ := λ (n : ℕ)
  (p : ∀ m : ℕ, n ≤ m → ℙ)
  (Hlen : p n (len n))
  (HleS : ∀ (m : ℕ) (H : n ≤ m),
    p (S m) (leS n m H)),
fix f (m : ℕ)(inst : n ≤ m)
  : p m inst :=
match inst as x
  return (p m x) with
  | len ⇒ Hlen
  | leS m x ⇒ HleS m x
end
  
```



AUXILIARY LEMMA GROUPS

- function properties and use cases
- basic lemmas
- typing of parts and types
- `wf_local` and `wf_local_rel` of environments
- equation for simplification



THE CASE CONSTRUCT

match m $t\text{Case } (\mathbb{N}, 0)$ type and
parameter count**as** n **in** P_n **with** $t\text{Lambda } "n"(t\text{Ind}\mathbb{N})$
 $(t\text{App } (t\text{Rel } 4) [t\text{Rel } 0])$
 $(t\text{Rel } 0)$

return type

match object

 $| 0 \Rightarrow P_0$ $(0, t\text{Rel } 2);$

branches

 $| S_n \Rightarrow P_{S_n}$ $(1, t\text{Lambda } "n"(t\text{Ind}\mathbb{N})$
 $(t\text{App } (t\text{Rel } 2) [t\text{Rel } 0]))$ **end**

THE CASE RULE

$$\begin{aligned}
 &\forall (\text{indnpar} : \text{inductive} \times \mathbb{N})(u : \text{universe_instance})(p\ c : \text{term}) \\
 &(\text{brs} : \mathcal{L}(\mathbb{N} \times \text{term}))(\text{args} : \mathcal{L}\ \text{term}), \\
 &\mathbf{let}\ \text{ind} := \text{indnpar}.1\ \mathbf{in} \\
 &\mathbf{let}\ \text{npar} := \text{indnpar}.2\ \mathbf{in} \\
 &\forall (\text{mdecl} : \text{mutual_inductive_body})(\text{idecl} : \text{one_inductive_body}), \\
 &\text{declared_inductive}\ \Sigma.1\ \text{mdecl}\ \text{idecl} \rightarrow \text{ind_npars}\ \text{mdecl} = \text{npar} \rightarrow \\
 &\mathbf{let}\ \text{params} := \text{firstn}\ \text{npar}\ \text{args}\ \mathbf{in}\ \forall (\text{ps} : \text{universe})(\text{pty} : \text{term}), \\
 &\text{build_case_predicate_type}\ \text{ind}\ \text{mdecl}\ \text{idecl}\ \text{params}\ u\ \text{ps} = \text{Some}\ \text{pty} \rightarrow \\
 &\Sigma; \Gamma \vdash p : \text{pty} \rightarrow \\
 &\text{existsb}\ (\text{leb_sort_family}\ (\text{universe_family}\ \text{ps}))(\text{ind_kelim}\ \text{idecl}) \rightarrow \\
 &\Sigma; \Gamma \vdash c : \text{mkApps}\ (\text{tInd}\ \text{ind}\ u)\ \text{args} \rightarrow \\
 &\forall\ \text{btys} : \mathcal{L}(\mathbb{N} \times \text{term}), \\
 &\text{map_option_out}\ (\text{build_branches_type}\ \text{ind}\ \text{mdecl}\ \text{idecl}\ \text{params}\ u\ \text{p}) = \text{Some}\ \text{btys} \rightarrow \\
 &\text{All2}\ (\lambda\ \text{br}\ \text{bty} : \mathbb{N} \times \text{term} \Rightarrow (\text{br}.1 = \text{bty}.1 \times \Sigma; \Gamma \vdash \text{br}.2 : \text{bty}.2)) \times \\
 &\Sigma; \Gamma \vdash \text{bty}.2 : \text{tSort}\ \text{ps})\ \text{brs}\ \text{btys} \rightarrow \\
 &\Sigma; \Gamma \vdash \text{tCase}\ \text{indnpar}\ p\ c\ \text{brs} : \text{mkApps}\ p\ (\text{skipn}\ \text{npar}\ \text{args}\ ++\ [c])
 \end{aligned}$$


DIFFICULTIES

- confusing lifting with indices
- search feature often does not work
- problems with unification
- no overview over the goal



ASSUMPTION

$$\Sigma; \Gamma, \text{params}, p, \text{cases}, \uparrow^{1+|\text{ctors}|} \text{indices} \vdash$$
$$t^{\text{ind_type}}$$
$$(\uparrow_{|\text{indices}|}^{1+|\text{ctors}|} \uparrow^{|\text{indices}|} \text{mkRel params} ++$$
$$\uparrow_{|\text{indices}|}^{1+|\text{ctors}|} \text{mkRel indices}) :$$
$$t\text{Sort } x$$


GOAL

$$\Sigma; \Gamma, \text{params}, p, \text{cases}, f, \uparrow^{2+|\text{ctors}|} \text{indices} \vdash$$

$$\uparrow t^{\text{ind_type}}$$

$$\left(\uparrow_{|\text{indices}|}^{2+|\text{ctors}|+|\text{indices}|} \text{mkRel params} ++ \right.$$

$$\left. \uparrow_{|\text{indices}|}^{1+|\text{ctors}|} \text{mkRel indices} \right):$$

$$\text{tSort } ?s$$


IDEA

$$\begin{aligned} & \Sigma; \Gamma, \mathbf{x}s \vdash t^a \mathbf{x} : \mathbb{T} \rightarrow \\ & \Sigma; \Gamma, \mathbf{y}s, \uparrow^{|\mathbf{y}s|} \mathbf{x}s \vdash \\ & t^{\uparrow^{|\mathbf{y}s|}_{|\mathbf{x}s|} a} \left(\uparrow^{|\mathbf{y}s|}_{|\mathbf{x}s|} \mathbf{x} \right) : \uparrow^{|\mathbf{y}s|}_{|\mathbf{x}s|} \mathbb{T} \end{aligned}$$

```
apply typing_spine_lifting.  
apply t0.
```



REALITY

```

1  instantiate (1 := x ).
   rewrite appLen, lift_context_len.
   cbn.
   rewrite appLen, revLen, quantifyCasesLen.
   rewrite mapLen.
   cbn.
   rewrite ind_arity_eq.
   rewrite removeArityCount.
2: eapply indParamVass;eassumption.
2: now rewrite uniP, nparCount.
   rewrite collectProdMkProd.
2: apply indicesVass.
2: cbn;congruence.
   rewrite revLen.
   unfold snoc, app_context.
   unfold snoc, app_context in t0.

rewrite ← appCons.
rewrite ← liftContextSucc.
set (xs:=lift_context ... ).
replace ((lift0 ... )
  (subst_instance_constr univ ... )) with
  (lift 1 #|xs| (subst_instance_constr ... )).
2: {
  subst xs.
  apply liftSubstInstanceConstr2.
}
change (tSort x ) with (lift 1 #|xs| (tSort x )).
replace (mapi

```

```

(λ (i : ℕ) _ =>
  tRel (... ))
  (ind_params mind_body ) ++
  map (lift (... ) (#|ind_indices| ))
  (nat_rec [] (λ (n : ℕ) (a : ℒ term) => tRel n :: a) #|ind
    with
      (map (lift 1 #|xs| )
        (mapi
          (λ (i : ℕ) _ =>
            lift ... #|ind_indices|
              ((lift0 ... )(tRel ... )))
            (ind_params mind_body ) ++
            map (lift ... #|ind_indices| )
              (nat_rec [] (λ (n : ℕ) (a : ℒ term) =>
                tRel (0 + n) :: a)
                #|ind_indices| )))
        45 2: { ... }
        69 subst xs.
           pose proof typingSpineLifting.
           unfold app_context in X.
           apply X.
           2: reflexivity.
           rewrite uniP, nparCount, sub_diag.
           unfold skipn.
           rewrite replaceUnderItMkProd.
           2: apply indicesVass.
           rewrite nparCount in t0.
           80 apply t0.

```



METACOQ DIFFICULTIES

- deeply nested definitions
- notational inconvenience
- function definitions with `fold_left`, `reverse` and `mapi`
- properties hidden in definitions and lemmas
- missing lemmas for properties of functions
- some lemmas and functions only in pCUIC

FUTURE WORK

- Move plugin to pCUIC
- Induction principles for mutual inductive types
- Stronger induction principles for nested inductive types

Inductive roseTree := tree (xs:list roseTree).



Scheme Induction **for** roseTree **Sort** T.

$\forall P : \text{roseTree} \rightarrow \mathbb{P},$
 $(\forall xs : \mathcal{L} \text{ roseTree}, P (\text{tree } xs)) \rightarrow$
 $\forall r : \text{roseTree}, P r$

MetaCoq Run Scheme Induction **for** roseTree **Sort** T.

$\forall P : \text{roseTree} \rightarrow \mathbb{P},$
 $(\forall xs : \mathcal{L} \text{ roseTree}, (\forall t, \text{In } t \text{ } xs \rightarrow P t) \rightarrow P (\text{tree } xs)) \rightarrow$
 $\forall r : \text{roseTree}, P r$

REFERENCES I

-  Abhishek Anand, Simon Boulrier, Cyril Cohen, Matthieu Sozeau, and Nicolas Tabareau, *Towards certified meta-programming with typed Template-Coq*, International Conference on Interactive Theorem Proving, Springer, 2018.
-  Matthieu Sozeau, Abhishek Anand, Simon Boulrier, Cyril Cohen, Yannick Forster, Fabian Kunze, Gregory Malecha, Nicolas Tabareau, and Théo Winterhalter, *The MetaCoq Project*.

INDUCTION

```

λ (A :  $\mathbb{T}$ ) (R : A → A →  $\mathbb{P}$ ) (p :  $\forall x : A, \text{Acc } R \ x \rightarrow \mathbb{T}$ )
(HAi :  $\forall (x : A) (H : \forall y : A, R \ y \ x \rightarrow \text{Acc } R \ y)$ ,
  ( $\forall (y : A) (H_0 : R \ y \ x), p \ y \ (H \ y \ H_0)$ ) →
  p x (Acc_intro x H)),
fix f (x : A) (inst : Acc R x) : p x inst :=
match inst as inst0 return (p x inst0) with
| Acc_intro x0 ⇒ HAi x x0
  ( $\lambda (y : A) (H : R \ y \ x), f \ y \ (x_0 \ y \ H)$ )
end

```

1. find arguments with inductive call as argument
2. extract indices
3. add inductive hypothesis
4. lift everything accordingly

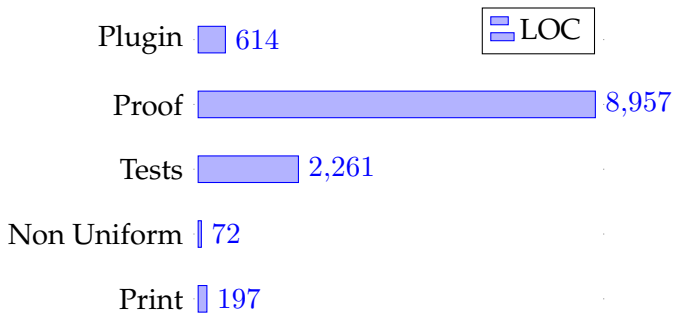


USED AXIOMS

- parameters and indices have empty bodies
- fix guard axiom
- induction properties



LINES OF CODE



GRAPH

